

Decomposition width—a new width parameter for matroids

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Abstract

We introduce a new width parameter for matroids called decomposition width and prove that every matroid property expressible in the monadic second order logic can be computed in linear time for matroids with bounded decomposition width if their decomposition is given. Since decompositions of small width for our new notion can be computed in polynomial time for matroids of bounded branch-width represented over finite fields, our results include recent algorithmic results of Hliněný [J. Combin. Theory Ser. B 96 (2006), 325–351] in this area and extend his results to matroids not necessarily representable over finite fields.

1 Introduction

Algorithmic aspects of graph tree-width form an important part of algorithmic graph theory. Efficient algorithms for computing tree-decompositions of graphs have been designed [1, 4] and a lot of NP-complete problems become tractable for classes of graphs with bounded tree-width [2, 3]. Most of results of the latter kind are implied by a general result of Courcelle [5, 6] that every graph property expressible in the monadic second-order logic can be decided in linear time for graphs with bounded tree-width.

As matroids are combinatorial structures that generalize graphs, it is natural to ask which of these results translate to matroids. Similarly, as in the case of graphs, some hard problems (that cannot be solved in polynomial time

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for general matroids) can be efficiently solved for (represented) matroids of small width. Though the notion of tree-width generalizes to matroids [16,17], a more natural width parameter for matroids is the notion of branch-width. Let us postpone a formal definition of this width parameter to Section 2 and just mention at this point that the branch-width of matroids is linearly related with their tree-width, in particular, the branch-width of a graphic matroid is bounded by twice the tree-width of the corresponding graph.

There are two main algorithmic aspects which one needs to address with respect to a width parameter of a combinatorial structure:

- the efficiency of computing decompositions of small width, and
- tractability of hard problems for input structures of small width.

The first issue has been successfully settled with respect to matroid branch-width: for every k , there exists an algorithm that either computes a branch-decomposition of an input matroid with width at most k or outputs that there is no such branch-decomposition. The first such algorithm has been found by Oum and Seymour [20] (an approximation algorithm has been known earlier [19], also see [10] for the case of matroids represented over a finite field) and a fixed parameter algorithm was later designed by Hliněný and Oum [15].

The tractability results, which include deciding monadic second-order logic properties [11,12], computing and evaluating the Tutte polynomial [14] and computing and counting representations over finite fields [18], usually require restricting to matroids represented over finite fields (see Section 2 for the definition). This is consistent with the facts that no subexponential algorithm can decide whether a given matroid is binary [22], i.e., representable over $\mathbb{GF}(2)$, even for matroids with bounded branch-width and that it is NP-hard to decide representability over $\mathbb{GF}(q)$ for every prime power $q \geq 4$ even for matroids with bounded branch-width represented over \mathbb{Q} [13], as well as with structural results on matroids [7–9] that also suggest that matroids representable over finite fields are close to graphic matroids (and thus graphs) but general matroids can be quite different.

The aim of this paper is to introduce another width parameter for matroids which will allow to extend the tractability results to matroids not necessarily representable over finite fields. The cost that needs to be paid for this is that this new width parameter cannot be bounded by the branch-width in general though it is bounded by the branch-width for matroids representable over a fixed finite field. Hence, we can obtain all tractability results mentioned in the previous paragraph. The new notion captures the

“structural finiteness” on cuts represented in the branch decomposition essential for the tractability results and is closely related to rooted configurations as introduced in [9] and indistinguishable sets from [18].

Let us state our results precisely. In Section 2, we introduce a K -decomposition of a matroid (where K is an integer) and define the decomposition width of a matroid \mathcal{M} to be the smallest integer K such that \mathcal{M} has a K -decomposition. In Section 3, we show that for every k and q , the decomposition width of a matroid with branch-width at most k that is representable over $\mathbb{GF}(q)$ is at most $F(k, q)$ and that $F(k, q)$ -decomposition of any such matroid can be computed in polynomial time. In Sections 5, 6 and 7, we show that for every K , there exist polynomial-time algorithms (with the degree of the polynomial independent of K) for computing and evaluating the Tutte polynomial, deciding monadic second-order logic properties, deciding representability and constructing and counting representations over finite fields when the input matroid is given by its K -decomposition. In particular, our results imply all the tractability results known for matroids represented over finite field (we here claim only the polynomiality, not matching the running times which we did not try to optimize throughout the paper).

A K -decomposition of a matroid actually captures the whole structure of a matroid, i.e., the matroid is fully described by its K -decomposition, and thus it can be understood as an alternative way of providing the input matroid. In fact, for a fixed K , the size of a K -decomposition is linear in the number of matroid elements and thus this representation of matroids is very suitable for this purpose. Let us state (without a proof) that K -decompositions of matroids *supports* contraction and deletion of matroid elements without increasing the decomposition width as well as some other matroid operations with increasing the decomposition width by a constant, e.g., relaxing a circuit-hyperplane increases the decomposition width by at most one. By *supporting* we mean that a K -decomposition of the new matroid can be efficiently computed from the K -decomposition of the original one. Hence, the definition of the decomposition width does not only yield a framework extending tractability results for matroids represented over finite fields with bounded branch-width but it also yields a compact data structure for representing input matroids.

2 Notation

In this section, we formally introduce the notions which are used in this paper. We start with basic notions and we then introduce matroid representations, branch-decompositions and our new width parameter. We also refer the

reader to the monographs [21, 23] for further exposition on matroids.

A matroid \mathcal{M} is a pair (E, \mathcal{I}) where $\mathcal{I} \subseteq 2^E$. The elements of E are called elements of \mathcal{M} , E is the ground set of \mathcal{M} and the sets contained in \mathcal{I} are called independent sets. The set \mathcal{I} is required to contain the empty set, to be hereditary, i.e., for every $F \in \mathcal{I}$, \mathcal{I} must contain all subsets of F , and to satisfy the exchange axiom: if F and F' are two sets of \mathcal{I} such that $|F| < |F'|$, then there exists $x \in F'$ such that $F \cup \{x\} \in \mathcal{I}$. The rank of a set F , denoted by $r(F)$, is the size of the largest independent subset of F (it can be inferred from the exchange axiom that all inclusion-wise maximal independent subsets of F have the same size). In the rest, we often understand matroids as sets of elements equipped with a property of “being independent”. We use $r(\mathcal{M})$ for the rank of the ground set of E .

If F is a set of elements of \mathcal{M} , then $\mathcal{M} \setminus F$ is the matroid obtained from \mathcal{M} by deleting the elements of F , i.e., the elements of $\mathcal{M} \setminus F$ are those not contained in F and a subset F' of such elements is independent in the matroid $\mathcal{M} \setminus F$ if and only if F' is independent in \mathcal{M} . The matroid \mathcal{M}/F which is obtained by contraction of F is the following matroid: the elements of \mathcal{M}/F are those not contained in F and a subset F' of such elements is independent in \mathcal{M}/F if and only if $r(F \cup F') = |F| + |F'|$. A loop of \mathcal{M} is an element e of \mathcal{M} such that $r(\{e\}) = 0$ and a bridge is an element such that $r(\mathcal{M} \setminus \{e\}) = r(\mathcal{M}) - 1$. A separation (A, B) is a partition of the elements of \mathcal{M} into two disjoint sets and a separation is called a k -separation if $r(\mathcal{M}) - r(A) - r(B) \leq k - 1$.

2.1 Matroid representations

Matroids do not generalize only the notion of graphs but they also generalize the notion of linear independence of vectors. If \mathbb{F} is a (finite or infinite) field, a mapping $\varphi : E \rightarrow \mathbb{F}^d$ from the ground set E of \mathcal{M} to a d -dimensional vector space over \mathbb{F} is a representation of \mathcal{M} if a set $\{e_1, \dots, e_k\}$ of elements of \mathcal{M} is independent in \mathcal{M} if and only if $\varphi(e_1), \dots, \varphi(e_k)$ are linearly independent vectors in \mathbb{F}^d . For a subset E of the elements of \mathcal{M} , $\varphi(E)$ denotes the linear subspace of \mathbb{F}^d generated by the images of the elements of E . In particular, $\dim \varphi(E) = r(E)$. Two representations φ_1 and φ_2 of \mathcal{M} are isomorphic if there exists an isomorphism ψ of vector spaces $\varphi_1(\mathcal{M})$ and $\varphi_2(\mathcal{M})$ such that $\psi(\varphi_1(e))$ is a non-zero multiple of $\varphi_2(e)$ for every element e of \mathcal{M} .

We next introduce additional notation for vector spaces over a field \mathbb{F} . If U_1 and U_2 are two linear subspaces of a vector space over \mathbb{F} , $U_1 \cap U_2$ is the linear space formed by all the vectors lying in both U_1 and U_2 , and $\mathcal{L}(U_1 \cup U_2)$ is the linear space formed by all the linear combinations of the vectors of U_1 and U_2 , i.e., the linear hull of $U_1 \cup U_2$. Formally, $v \in \mathcal{L}(U_1 \cup U_2)$ if and only

if there exist $v_1 \in U_1$, $v_2 \in U_2$ and $\alpha_1, \alpha_2 \in \mathbb{F}$ such that $v = \alpha_1 v_1 + \alpha_2 v_2$.

2.2 Branch-decompositions

A branch-decomposition of a matroid \mathcal{M} with ground set E is a tree T such that

- all inner nodes of T have degree three, and
- the leaves of T one-to-one correspond to the elements of \mathcal{M} .

An edge e of T splits T into two subtrees and the elements corresponding to the leaves of the two subtrees form a partition (E_1, E_2) of the ground set E . The width of an edge e is equal to $r(E_1) + r(E_2) - r(E)$, i.e., to the smallest k such that (E_1, E_2) is a $(k + 1)$ -separation of \mathcal{M} . The width of the branch-decomposition T is the maximum width of an edge e of T . Finally, the branch-width of a matroid is the minimum width of a branch-decomposition of \mathcal{M} and is denoted by $\text{bw}(\mathcal{M})$.

2.3 Decomposition width

We now formally define our new width parameter. A K -decomposition, $K \geq 1$, of a matroid \mathcal{M} is a rooted tree \mathcal{T} such that

- the leaves of \mathcal{T} one-to-one correspond to the elements of \mathcal{M} and carry the information whether the associated element of \mathcal{M} is a loop, and
- each inner node v of \mathcal{T} has exactly two children and is associated with two functions φ_v and φ_v^r such that $\varphi_v : \{0, \dots, K\}^2 \rightarrow \{0, \dots, K\}$ and $\varphi_v^r : \{0, \dots, K\}^2 \rightarrow \mathbb{N}$.

The tree \mathcal{T} represents the rank function of \mathcal{M} in the way we now describe. For a subset F of the ground set of \mathcal{M} , we label and color the vertices of \mathcal{T} in the following way: a leaf of \mathcal{T} corresponding to an element in F is colored with 1 and other leaves are colored with 0. The leaves colored with 1 corresponding to non-loop elements of \mathcal{M} are labeled with 1 and all other leaves are labeled with 0. If v is an inner node of \mathcal{T} and its two children are labeled with λ_1 and λ_2 and colored with γ_1 and γ_2 , the node v is colored with $\varphi_v(\gamma_1, \gamma_2)$ and labeled with the number $\lambda_1 + \lambda_2 - \varphi_v^r(\gamma_1, \gamma_2)$. The rank of F is the label of the root of \mathcal{T} . The decomposition width of a matroid \mathcal{M} is the smallest K such that \mathcal{M} has a K -decomposition.

Let us give an intuitive explanation of the above procedure. The colors represent types of different subsets of elements of \mathcal{M} , i.e., if E_v is the set

of elements assigned to leaves of a subtree rooted at an inner vertex v , then those subsets of E_v that get the same color at v are of the same type. The labels represent the rank of subsets. Subsets of the same type can have different labels (and thus ranks) but they behave in the same way in the following sense: if E_1 and E_2 are elements assigned to leaves of two subtrees rooted at children of v , then the rank of the union of two subsets of E_1 with the same color and two subsets of E_2 with the same color is equal to the sum of their ranks decreased by the same amount.

In what follows, we will always assume that if F is the empty set, then all the nodes of \mathcal{T} are colored and labelled with 0 and consider this assumption to be part of the definition of a K -decomposition.

3 Constructing decompositions

In this section, we relate the decomposition width of matroids representable over finite fields to their branch-width.

Theorem 1. *Let \mathcal{M} be a matroid representable over a finite field \mathbb{F} of order q . If the branch-width of \mathcal{M} is at most $k \geq 1$, then the decomposition width of \mathcal{M} is at most $K = \frac{q^{k+1} - q(k+1) + k}{(q-1)^2}$. Moreover, if a branch-decomposition of \mathcal{M} with width k and its representation over \mathbb{F} are given, then a K -decomposition of \mathcal{M} can be constructed in time $O(n^{1+\alpha})$ where n is the number of elements of \mathcal{M} and α is the exponent from the matrix multiplication algorithm.*

Proof. Let \mathcal{T}_b be the branch-decomposition of \mathcal{M} of width at most k . Subdivide an arbitrary edge of \mathcal{T}_b and root the resulting tree \mathcal{T} at the new vertex. We now have to construct the functions φ_v and φ'_v . Fix a representation of \mathcal{M} over \mathbb{F} and let w_e be a vector representing an element e of \mathcal{M} . For a node v of \mathcal{T} , let W_v be the linear hull of all the vectors w_e representing the elements e corresponding to the leaves of the subtree rooted at v and W'_v be the linear hull of the remaining vectors w_e , i.e., vectors w_e for the elements e not corresponding to the leaves of the subtree. Finally, let D_v be the intersection of W_v and W'_v . Since \mathcal{T}_b is a branch-decomposition of width at most k , the dimension of the subspace $D_v = W_v \cap W'_v$ is at most k for every inner node v .

Each of at most $K + 1 = 1 + \sum_{i=1}^k \frac{q^i - 1}{q - 1}$ subspaces of D_v is associated with one of the colors $\{0, \dots, K\}$ in such a way that the trivial subspace of dimension 0 is associated with the color 0. If v is a leaf associated with non-trivial subspace D_v (the corresponding element is neither loop nor co-loop), then the subspace D_v itself is associated with the color 1. Let v be an inner node of \mathcal{T} with children v_1 and v_2 . For $\gamma_i \in \{0, \dots, K\}$, $i = 1, 2$, let Γ_i be

the subspace of W_i colored with γ_i . The function $\varphi_v(\gamma_1, \gamma_2)$ is equal to the color z of the subspace $D_v \cap \mathcal{L}(\Gamma_1 \cup \Gamma_2)$ where $\mathcal{L}(\Gamma_1 \cup \Gamma_2)$ is the linear hull of $\Gamma_1 \cup \Gamma_2$ and $\varphi_v^r(\gamma_1, \gamma_2)$ is equal

$$\dim \Gamma_1 + \dim \Gamma_2 - \dim \mathcal{L}(\Gamma_1 \cup \Gamma_2) .$$

If there is no subspace of D_{v_i} with color γ_i for $i = 1$ or $i = 2$, then the functions $\varphi_v(\gamma_1, \gamma_2)$ and $\varphi_v^r(\gamma_1, \gamma_2)$ are equal to 0. Finally, if v is the root of the tree \mathcal{T} , then the function φ_v^r is defined in the same way and φ_v is defined to be constantly equal to 0.

We now have to verify that the constructed decomposition of \mathcal{M} represents the matroid \mathcal{M} . Let F be a subset of elements of \mathcal{M} and consider the coloring and the labeling of the nodes of \mathcal{M} as in the definition of a K -decomposition. We will prove that the label of each node v of \mathcal{T} is equal to the dimension of the linear hull X_v of the vectors w_e with $e \in F_v$ where F_v is the set of the elements of F corresponding to the leaves of the subtree rooted at v .

The label of each leaf v is equal to the dimension of X_v since the leaves corresponding to elements e of F that are not loops are labeled with 1 and the other leaves are labeled with 0. Let v be a node of \mathcal{T} with children v_1 and v_2 . Let us compute the dimension of X_v :

$$\begin{aligned} \dim X_v &= \dim X_{v_1} + \dim X_{v_2} - \dim X_{v_1} \cap X_{v_2} \\ &= \dim X_{v_1} + \dim X_{v_2} - \dim X_{v_1} \cap X_{v_2} \cap D_{v_1} \cap D_{v_2} \\ &= \dim X_{v_1} + \dim X_{v_2} + \\ &\quad \dim X_{v_1} \cap D_{v_1} + \dim X_{v_2} \cap D_{v_2} - \\ &\quad \dim \mathcal{L}((X_{v_1} \cap D_{v_1}) \cup (X_{v_2} \cap D_{v_2})) \end{aligned}$$

since $X_{v_1} \cap X_{v_2} \subseteq W_{v_i} \cap W'_{v_i} = D_{v_i}$. As

$$\varphi_v^r(\gamma_1, \gamma_2) = \dim X_{v_1} \cap D_{v_1} + \dim X_{v_2} \cap D_{v_2} - \dim \mathcal{L}((X_{v_1} \cap D_{v_1}) \cup (X_{v_2} \cap D_{v_2}))$$

where γ_i is the color of $X_{v_i} \cap D_{v_i}$, the label of v is equal to the dimension of X_v . This finishes the proof that the decomposition \mathcal{T} represents the matroid \mathcal{M} .

It remains to explain how to construct the decomposition \mathcal{T} within the claimed time. The tree \mathcal{T} has $n - 1$ non-leaf nodes v and we need time $O(n^\alpha)$ to compute the subspaces W_v , W'_v and D_v . Since the dimension of D_v is bounded by k , associating colors $\{0, \dots, K\}$ with the subspaces of D_v and defining the functions φ_v and φ_v^r requires time bounded by K and q for each node v . \square

Theorem 1 and the cubic-time algorithm of Hliněný and Oum [15] for computing branch-decompositions of matroids with bounded branch-width yield the following:

Corollary 2. *Let \mathcal{M} be a matroid represented over a finite field \mathbb{F} of order q . For every $k \geq 1$, there exists an algorithm running in time $O(n^{1+\alpha})$, where n is the number of elements of \mathcal{M} and α is the exponent from the matrix multiplication algorithm, that either outputs that the branch-width of \mathcal{M} is bigger than k or construct a K -decomposition of \mathcal{M} for $K \leq \frac{q^{k+1}-q(k+1)+k}{(q-1)^2}$.*

4 Verifying decompositions

A K -decomposition fully describes a considered matroid \mathcal{M} through functions φ_v and φ_v^r for internal nodes v of the decomposition. Clearly, not all possible choices of φ_v and φ_v^r give rise to a decomposition representing a matroid and it is natural to ask whether we can efficiently test that a K -decomposition represents a matroid. We answer this equation in the affirmative way in the next theorem.

Theorem 3. *For every K -decomposition \mathcal{T} with n leaves, it can be tested in time $O(K^8 n)$ whether \mathcal{T} corresponds to a matroid.*

Proof. One of the equivalent definitions of a matroid is given by the submodularity of its rank function, i.e., a function $r : 2^E \rightarrow \mathbb{N}$ is a rank function of a matroid if and only if

$$r(A \cup B) + r(A \cap B) \leq r(A) + r(B) \tag{1}$$

for every subsets A and B of E . We will verify using the K -decomposition that the inequality (1) holds for all $A, B \subseteq E$.

We first check that for the empty set, all nodes of \mathcal{T} are colored with 0 and labelled with 0. If this is not the case, we reject. Let E_v be now the set of the elements assigned to the leaves of a subtree of \mathcal{T} rooted at v . In order to achieve our goal, we will compute for every node v of \mathcal{T} and every quadruple $[\gamma_A, \gamma_B, \gamma_\cap, \gamma_\cup] \in \{0, \dots, K\}^4$ the minimum possible difference

$$\lambda(A) + \lambda(B) - \lambda(A \cup B) - \lambda(A \cap B)$$

where A is a subset of E_v colored at v with γ_A and B a subset colored with γ_B such that the color of $A \cup B$ is γ_\cup and the color of $A \cap B$ is γ_\cap and $\lambda(A)$, $\lambda(B)$, $\lambda(A \cup B)$ and $\lambda(A \cap B)$ are the labels corresponding to the subsets A , B , $A \cup B$ and $A \cap B$ at v when the coloring and labelling procedure from the

definition of a K -decomposition is applied. This minimum is further denoted by $\mu_v(\gamma_A, \gamma_B, \gamma_\cap, \gamma_\cup)$. If there is no pair of subsets A and B such that the color of A at v is γ_A , the color of B is γ_B , the color of $A \cup B$ is γ_\cup and the color of $A \cap B$ is γ_\cap , then $\mu_v(\gamma_A, \gamma_B, \gamma_\cap, \gamma_\cup) = \infty$.

Let us see how the values of μ can be computed for the nodes of \mathcal{T} in the direction from the leaves towards the root of \mathcal{T} . If v is a leaf of \mathcal{T} , then the function μ_v is equal to zero for the quadruples $[0, 0, 0, 0]$, $[1, 0, 0, 1]$, $[0, 1, 0, 1]$ and $[1, 1, 1, 1]$ and is equal to ∞ for all other quadruples. If v is an inner node of \mathcal{T} with children v' and v'' , then $\mu_v(\gamma_A, \gamma_B, \gamma_\cap, \gamma_\cup)$ is equal to the minimum of

$$\begin{aligned} & \mu_{v'}(\gamma'_A, \gamma'_B, \gamma'_\cap, \gamma'_\cup) + \mu_{v''}(\gamma''_A, \gamma''_B, \gamma''_\cap, \gamma''_\cup) - \\ & \varphi_v^r(\gamma'_A, \gamma''_A) - \varphi_v^r(\gamma'_B, \gamma''_B) + \varphi_v^r(\gamma'_\cap, \gamma''_\cap) + \varphi_v^r(\gamma'_\cup, \gamma''_\cup) \end{aligned}$$

where the minimum is taken over all pairs of quadruples $[\gamma'_A, \gamma'_B, \gamma'_\cap, \gamma'_\cup]$ and $[\gamma''_A, \gamma''_B, \gamma''_\cap, \gamma''_\cup]$ such that $\varphi_v(\gamma'_A, \gamma''_A) = \gamma_A$, $\varphi_v(\gamma'_B, \gamma''_B) = \gamma_B$, $\varphi_v(\gamma'_\cap, \gamma''_\cap) = \gamma_\cap$ and $\varphi_v(\gamma'_\cup, \gamma''_\cup) = \gamma_\cup$. Here, we assume that $\infty + k = \infty$ and $k < \infty$ for any $k \in \mathbb{N}$. It is straightforward to verify that the function μ_v computed in this way is equal to the claimed minimum.

The function μ_v for the root v of \mathcal{T} determines whether the K -decomposition represents a matroid. If the value of μ_v for any quadruple is negative, then there exists subsets A and B violating (1) and thus the rank function given by \mathcal{T} is not submodular. On the other hand, if all the values of μ_v are non-negative or equal to ∞ for all quadruples, then (1) holds for every two subsets A and B . Consequently, \mathcal{T} represents a matroid.

It remains to estimate the running time of the algorithm. Clearly, the minimums can be computed in time $O(K^8)$ for every inner node of \mathcal{T} . We conclude that the running time of the whole algorithm is $O(K^8 n)$ as claimed. \square

5 Computing the Tutte polynomial

One of the classical polynomials associated to matroids is the Tutte polynomial. There are several equivalent definitions of this polynomial but we provide here only the one we use. For a matroid \mathcal{M} with ground set E , the Tutte polynomial $T_{\mathcal{M}}(x, y)$ is equal to

$$T_{\mathcal{M}}(x, y) = \sum_{F \subseteq E} (x - 1)^{r(E) - r(F)} (y - 1)^{|F| - r(F)} \quad (2)$$

The Tutte polynomial is an important algebraic object associated to a matroid. Some of the values of $T_{\mathcal{M}}(x, y)$ have a combinatorial interpretation;

as a simple example, the value $T_{\mathcal{M}}(1, 1)$ is equal to the number of bases of a matroid \mathcal{M} . We show that the Tutte polynomial can be computed and evaluated in time $O(K^2n^3r^2)$ for n -element matroids of rank r given by the K -decomposition. The part of the proof for computing the Tutte polynomial reflects the main motivation behind the definition of matroid decompositions given in the previous section.

Theorem 4. *Let K be a fixed integer. The Tutte polynomial of an n -element matroid \mathcal{M} given by its K -decomposition can be computed in time $O(K^2n^3r^2)$ and evaluated in time $O(K^2n)$ where r is the rank of \mathcal{M} (under the assumption that summing and multiplying $O(n)$ -bit numbers can be done in constant time).*

Proof. First, we give an algorithm for computing the Tutte polynomial, i.e., an algorithm that computes the coefficients in the polynomial. Let \mathcal{T} be a K -decomposition of \mathcal{M} and let E_v be the elements of \mathcal{M} corresponding to leaves of the subtree rooted at a vertex v . For every node v of \mathcal{M} and every triple $[\gamma, n', r']$, $0 \leq \gamma \leq K$, $0 \leq n' \leq n$ and $0 \leq r' \leq r$, we compute the number of subsets F of E_v such that the color assigned to v for F is γ , $|F| = n'$ and the rank of F is r' . These numbers will be denoted by $\mu_v(\gamma, n', r')$.

We compute the numbers $\mu_v(\gamma, n', r')$ from the leaves to the root of \mathcal{T} . If v is a leaf of \mathcal{T} , then $\mu_v(0, 1, 0)$ is equal to 2 if v corresponds to a loop of \mathcal{M} and it is equal to 1, otherwise. In the latter case, i.e., if v corresponds to a non-loop, $\mu_v(1, 1, 1)$ is also equal to 1. All the other values of μ_v are equal to 0.

Let v be a node with children v_1 and v_2 . If F_i is a subset of E_{v_i} with color γ_i such that $n_i = |F_i|$ and $r_i = r(F_i)$, then $F_1 \cup F_2$ is a subset of E_v with color $\varphi_v(\gamma_1, \gamma_2)$, with $n_1 + n_2$ elements and the rank $r_1 + r_2 - \varphi_v^r(\gamma_1, \gamma_2)$. In other words, it holds that

$$\mu_v(\gamma, n', r') = \sum_{\gamma_1, n_1, r_1, \gamma_2, n_2, r_2} \mu_{v_1}(\gamma_1, n_1, r_1) \mu_{v_2}(\gamma_2, n_2, r_2) \quad (3)$$

where the sum is taken over six-tuples $(\gamma_1, n_1, r_1, \gamma_2, n_2, r_2)$ such that $\gamma = \varphi_v(\gamma_1, \gamma_2)$, $n' = n_1 + n_2$ and $r' = r_1 + r_2 - \varphi_v^r(\gamma_1, \gamma_2)$. Computing μ_v from the values of μ_{v_1} and μ_{v_2} base on (3) requires time $O(K^2n^2r^2)$ and thus the total running time of the algorithm is $O(K^2n^3r^2)$. The Tutte polynomial of \mathcal{M} can be read from μ_r where r is the root of \mathcal{T} since the value $\mu_r(0, \alpha, \beta)$ is the coefficient at $(x-1)^{r(E)-\beta}(y-1)^{\alpha-\beta}$ in (2).

Let us turn our attention to evaluating the Tutte polynomial for given values of x and y . This time, we recursively compute the following quantity

for every node v of \mathcal{T} :

$$\mu_v(\gamma) = \sum_{F \subseteq E_v} (x-1)^{r(E)-r(F)} (y-1)^{|F|-r(F)} \quad (4)$$

where the sum is taken over the subsets F with color γ . The value of $\mu_v(0)$ is equal to $(x-1)^{r(E)}y$ for a leaf v corresponding to a loop of \mathcal{M} and to $(x-1)^{r(E)}$ if v correspond to a non-loop. In the latter case, $\mu_v(1)$ is also equal to $(x-1)^{r(E)-1}$. All other values of μ_v are equal to 0 in both cases.

For a node v of \mathcal{T} with two children v_1 and v_2 , the equation (4) and the definition of a K -decomposition implies that

$$\mu_v(\gamma) = \sum_{0 \leq \gamma_1, \gamma_2 \leq K} \frac{\mu_{v_1}(\gamma_1) \mu_{v_2}(\gamma_2)}{(x-1)^{r(E)+\varphi_v^r(\gamma_1, \gamma_2)} (y-1)^{\varphi_v^r(\gamma_1, \gamma_2)}}$$

where the sum is taken over values γ_1 and γ_2 such that $\gamma = \varphi_v(\gamma_1, \gamma_2)$. Under the assumption that arithmetic operations require constant time, determining the values of μ_v needs time $O(K^2)$. Since the number of nodes of \mathcal{T} is $O(n)$, the total running time of the algorithm is $O(K^2n)$ as claimed in the statement of the theorem. \square

As a corollary, we obtain Hliněný's result on computing the Tutte polynomial and its values for matroids represented over finite fields of bounded branch-width [14].

Corollary 5. *Let \mathbb{F} be a fixed finite field and k a fixed integer. There is a polynomial-time algorithm for computing and evaluating the Tutte polynomial for the class of matroids of branch-width at most k representable over \mathbb{F} that are given by their representation over the field \mathbb{F} .*

Proof. Let \mathcal{M} be a matroid of branch-width at most k represented over \mathbb{F} . By Corollary 2, we can efficiently find a K -decomposition of \mathcal{M} where K is a constant depending only on \mathbb{F} and k . As K depends on \mathbb{F} and k only, the Tutte polynomial of \mathcal{M} can be computed and evaluated in time polynomial in the number of elements of \mathcal{M} by Theorem 4. \square

6 Deciding MSOL-properties

In this section, we show that there is a linear-time algorithm for deciding monadic second order logic formulas for matroids of bounded decomposition width. First, let us be precise on the type of formulas that we are interested

in. A monadic second order logic formula ψ , an MSOL formula, for a matroid contains basic logic operators (the negation, the disjunction and the conjunction), quantifications over elements and sets of elements of a matroid (we refer to these variables as to element and set variables), the equality predicate, the predicate of containment of an element in a set and the independence predicate which determines whether a set of elements of a matroid is independent. The independence predicate depends on and encodes an input matroid \mathcal{M} .

In order to present the algorithm, we introduce auxiliary notions of a K -half and K -halved matroids which we extend to interpreted K -halves. A K -half is a matroid \mathcal{M} with ground-set E and equipped with a function $\varphi_{\mathcal{M}} : 2^E \rightarrow \{0, \dots, K\}$. A K -halved matroid is a matroid \mathcal{M} with ground-set $E = E_1 \cup E_2$ composed of a K -half \mathcal{M}_1 with ground set E_1 and a matroid \mathcal{M}_2 with ground set E_2 such that each subset of $F \subseteq E_2$ is assigned a vector w^F of $K + 1$ non-negative integers. Both \mathcal{M}_1 and \mathcal{M}_2 are matroids. The rank of a subset $F \subseteq E_1 \cup E_2$ is given by the following formula:

$$r_{\mathcal{M}}(F) = r_{\mathcal{M}_1}(F \cap E_1) + r_{\mathcal{M}_2}(F \cap E_2) - w_{\varphi_{\mathcal{M}_1}(F \cap E_1)}^{F \cap E_2},$$

where $w_{\varphi_{\mathcal{M}_1}(F \cap E_1)}^{F \cap E_2}$ is the coordinate of the vector $w^{F \cap E_2}$ corresponding to $\varphi_{\mathcal{M}_1}(F \cap E_1)$. We will write $\mathcal{M} = \mathcal{M}_1 \oplus_K \mathcal{M}_2$ to represent the fact that the matroid \mathcal{M} is a K -halved matroid obtained from \mathcal{M}_1 and \mathcal{M}_2 in the way we have just described.

The next lemma justifies the just introduced definition of K -halved matroids since it asserts that every matroid \mathcal{M} represented by a K -decomposition can be viewed as a composed of two K -halves, one of which is \mathcal{M} restricted to the elements corresponding to leaves of a subtree of its K -decomposition.

Lemma 6. *Let \mathcal{T} be a K -decomposition of a matroid \mathcal{M} , $K \geq 1$, and let v be a node of \mathcal{T} . Further, let E_v be the set of elements of \mathcal{M} assigned to the leaves of the subtree of \mathcal{T} rooted at v , and $\overline{E_v}$ the set of the remaining elements of \mathcal{M} . If F_1 and F_2 are two subsets of E_v such that the color of v with respect to the decomposition is the same for F_1 and F_2 , then*

$$r(F) + r(F_1) - r(F \cup F_1) = r(F) + r(F_2) - r(F \cup F_2)$$

for every subset F of $\overline{E_v}$.

Proof. The equality

$$r(F) + r(F_1) - r(F \cup F_1) = r(F) + r(F_2) - r(F \cup F_2)$$

follows from the fact that the vertex v gets the same color for both sets F_1 and F_2 : switching between F_1 and F_2 can change the labels of the vertices in the subtree of v and on the path P from v to the root only. Since the color of v is the same for F_1 and F_2 , switching between F_1 and F_2 does not change the colors of the vertices on the path P . Hence, the label of the root is equal to the sum of the labels of the children of the vertices on P decreased by the values of the functions φ_w for vertices on P which are not affected by switching between F_1 and F_2 since they only depend on the colors of children of w which are not changed by the switch. Hence, the equality from the statement of the lemma holds. \square

For the purpose of induction in the next lemma, let us allow free variables in formulas. For an MSOL formula ψ with free variables ξ_1, \dots, ξ_k , a K -signature σ is a mapping $\sigma : \{1, \dots, k\} \rightarrow \{0, \dots, K\} \cup \{\star\}$ such that $\sigma^{-1}(\star)$ contains only indices of element variables. A σ -interpretation of a K -half \mathcal{M}_1 is an assignment of elements e to element variables ξ_j with $\sigma(j) \neq \star$ such that $\varphi_{\mathcal{M}_1}(\{e\}) = \sigma(j)$ and an assignment of subsets F of the ground set of \mathcal{M}_1 to set variables ξ_j with $\varphi_{\mathcal{M}_1}(F) = \sigma(j)$. A σ -interpretation of \mathcal{M}_2 is an assignment of elements of \mathcal{M}_2 to element variables ξ_j with $\sigma(j) = \star$ and an assignment of subsets F of the ground set of \mathcal{M}_2 to set variables ξ_j . A matroid \mathcal{M}_i , $i \in \{1, 2\}$, with a σ -interpretation is said to be σ -interpreted.

If both \mathcal{M}_1 and \mathcal{M}_2 are σ -interpreted, then the σ -interpretations of \mathcal{M}_1 and \mathcal{M}_2 naturally give an assignment of element variables ξ_j (given by the assignment for \mathcal{M}_1 if $\sigma(j) \neq \star$ and by the assignment for \mathcal{M}_2 , otherwise) and an assignment of set variables ξ_j by uniting the assignments for ξ_j in the σ -interpretations of \mathcal{M}_1 and \mathcal{M}_2 . For a formula ψ with a K -signature σ , ψ is satisfied for a matroid $\mathcal{M} = \mathcal{M}_1 \oplus_K \mathcal{M}_2$ with a σ -interpreted K -half \mathcal{M}_1 and σ -interpreted \mathcal{M}_2 if the assignment given by the σ -interpretations and \mathcal{M} satisfies ψ .

The crucial notion in our argument is that of (ψ, σ) -equivalence of σ -interpreted K -halves which we now define. For a formula ψ with a K -signature σ , two σ -interpreted K -halves \mathcal{M}_1 and \mathcal{M}'_1 are (ψ, σ) -equivalent if for every σ -interpreted \mathcal{M}_2 , the formula ψ is satisfied for $\mathcal{M}_1 \oplus_K \mathcal{M}_2$ if and only if it is satisfied for $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$.

In the next lemma, we show that the number of (ψ, σ) -equivalence classes of σ -interpreted K -halves is finite for every MSOL formula ψ and every K -signature σ of ψ .

Lemma 7. *Let ψ be a fixed MSOL formula and σ a K -signature of ψ . The number of (ψ, σ) -equivalence classes of K -halves is finite.*

Proof. We prove the statement of the lemma by induction on the size of

ψ . Let us start with the simplest possible formulas, i.e., $\xi_1 = \xi_2$, $\xi_1 \in \xi_2$ and $\text{ind}_{\mathcal{M}}(\xi_1)$ where $\text{ind}_{\mathcal{M}}$ is the independence predicate for a matroid $\mathcal{M} = \mathcal{M}_1 \oplus_K \mathcal{M}_2$.

Let ψ be equal to $\xi_1 = \xi_2$. If $\sigma(1) \neq \sigma(2)$, then ψ cannot be σ -satisfied and thus all σ -interpreted K -halves are (ψ, σ) -equivalent. If $\sigma(1) = \sigma(2) \neq \star$ and ξ_1 and ξ_2 are element variables, then two σ -interpreted K -halves \mathcal{M}_1 and \mathcal{M}'_1 are (ψ, σ) -equivalent if their interpretations assign ξ_1 and ξ_2 the same element. Hence, the number of (ψ, σ) -equivalence classes of σ -interpreted K -halves is two. If $\sigma(1) = \sigma(2) = \star$, then all σ -interpreted K -halves are (ψ, σ) -equivalent since the truth of ψ is determined by the σ -interpretation of \mathcal{M}_2 . If ξ_1 and ξ_2 are set variables, then a little more complex argument shows that the number of (ψ, σ) -equivalence classes of σ -interpreted K -halves is two if $\sigma(1) = \sigma(2)$ and it is one if $\sigma(1) \neq \sigma(2)$.

For the formula $\psi = (\xi_1 \in \xi_2)$, the number of equivalence classes for $\sigma(1) = \star$ is one as the truth of the formula ψ depends only on the σ -interpretation of \mathcal{M}_2 and the number of (ψ, σ) -equivalence classes for $\sigma(1) \neq \star$ is two depending on the fact whether $\xi_1 \in \xi_2$ in the σ -interpretation of \mathcal{M}_1 . For the formula $\psi = \text{ind}_{\mathcal{M}}(\xi_1)$, there are two (ψ, σ) -equivalence classes distinguishing σ -interpreted K -halves with sets ξ_1 independent in the σ -interpretation and those with a set ξ_1 that is not independent.

We now consider more complex MSOL formulas. By standard logic manipulation, we can assume that ψ is of one of the following forms: $\neg\psi_1$, $\psi_1 \vee \psi_2$ or $\exists\xi\psi_1$. The (ψ, σ) -equivalence classes for $\psi = \neg\psi_1$ are the same as the (ψ_1, σ) -equivalence classes. Hence, their number is finite by the induction.

For the disjunction $\psi_1 \vee \psi_2$, if \mathcal{M}_1 and \mathcal{M}'_1 are σ -interpreted K -halves that are both (ψ_1, σ) -equivalent and (ψ_2, σ) -equivalent, then \mathcal{M}_1 and \mathcal{M}'_1 are also (ψ, σ) -equivalent. Indeed, if $\mathcal{M}_1 \oplus_K \mathcal{M}_2$ satisfies ψ_j , $j \in \{1, 2\}$, then $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$ also satisfies ψ_j . On the other hand, if $\mathcal{M}_1 \oplus_K \mathcal{M}_2$ satisfies neither ψ_1 nor ψ_2 , then $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$ also satisfies neither ψ_1 nor ψ_2 . Hence, the (ψ, σ) -equivalence classes of σ -interpreted K -halves are unions of the intersections of the (ψ_1, σ) -equivalence classes and (ψ_2, σ) -equivalence classes. In particular, the number of (ψ, σ) -equivalence classes of σ -interpreted K -halves is finite.

The final type of formulas we consider are those of the type $\psi = \exists\xi\psi_1$. Let \mathcal{M}_1 and \mathcal{M}'_1 be two σ -interpreted K -halves. Note that the domain of σ' contains in addition to σ the index corresponding to the variable ξ . There are $K+1$ or $K+2$ possible extensions of σ to σ' (the number depends on the fact whether ξ is an element or set variable). Let $\sigma'_1, \dots, \sigma'_N$, $N \in \{K+1, K+2\}$, be the extensions of σ and let C_i , $i = 1, \dots, N$, be the number of (ψ, σ'_i) -equivalence classes of σ'_i -interpreted K -halves. Every σ -interpretation of \mathcal{M}_1 can be extended to a σ'_i -interpretation of \mathcal{M}_1 in possibly C_i non-equivalent ways.

We claim that if the extensions of the σ -interpretation of \mathcal{M}_1 and \mathcal{M}'_1 to σ'_i -interpretations appear in exactly the same (ψ, σ'_i) -equivalence classes for every i , then \mathcal{M}_1 and \mathcal{M}'_1 are (ψ, σ) -equivalent. Since there are $2^{C_1+\dots+C_N}$ possible types of extensions of the σ -interpretation to σ'_i -interpretations and those \mathcal{M}_1 and \mathcal{M}'_1 with the extensions of the same type must be (ψ, σ) -equivalent, the number of (ψ, σ) -equivalence classes is bounded by $2^{C_1+\dots+C_N}$.

Assume that \mathcal{M}_1 and \mathcal{M}'_1 with extensions of the same type are not (ψ, σ) -equivalent. By symmetry, we can assume that there exist a σ -interpreted \mathcal{M}_2 such that ψ is satisfied in $\mathcal{M}_1 \oplus_K \mathcal{M}_2$ and is not satisfied in $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$. Let ξ_0 be the choice of ξ that satisfies ψ_1 in $\mathcal{M}_1 \oplus_K \mathcal{M}_2$. The value ξ_0 uniquely determines an extension of σ to σ'_i by setting the value of σ'_i corresponding to ξ to $\varphi_{\mathcal{M}_1}(\xi_0)$ or $\varphi_{\mathcal{M}_1}(\{\xi_0\})$ unless ξ_0 is an element of \mathcal{M}_2 ; if ξ_0 is an element of \mathcal{M}_2 , the value of σ'_i corresponding to ξ should be set to \star . Since ψ_1 is satisfied for the σ -interpretations of \mathcal{M}_1 and \mathcal{M}_2 by choosing $\xi = \xi_0$, ψ_1 is satisfied for a σ'_i -interpretation of \mathcal{M}_1 and a σ'_i -interpretation of \mathcal{M}_2 . By our assumption, there exists a σ'_i -interpretation of \mathcal{M}'_1 that extends its σ -interpretation and that is (ψ, σ'_i) -equivalent to the σ'_i -interpretation of \mathcal{M}_1 . By the definition of (ψ, σ'_i) -equivalence, the formula ψ is satisfied in $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$. In particular, there exists a choice of the value of ξ in $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$ such that $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$ is satisfied for the σ -interpretations of \mathcal{M}'_1 and \mathcal{M}_2 which contradicts our assumption that ψ is not satisfied in $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$. \square

Let us return to our original problem of deciding MSOL formulas with no free variables. Analogously, for an MSOL formula ψ with no free variables, two K -halves \mathcal{M}_1 and \mathcal{M}'_1 are ψ -equivalent if the formula ψ is satisfied for $\mathcal{M}_1 \oplus_K \mathcal{M}_2$ if and only if ψ is satisfied for $\mathcal{M}'_1 \oplus_K \mathcal{M}_2$. By Lemma 7, the number of ψ -equivalence classes of K -halves is finite.

For a K -decomposition \mathcal{T} of a matroid \mathcal{M} , we can obtain a K -half by restricting \mathcal{M} to the elements corresponding to the leaves of a subtree of \mathcal{T} (note that the subsets of the elements not corresponding to the leaves of \mathcal{T} can be assigned non-negative integers as in the definition of a K -halved matroid by Lemma 6). The K -half obtained from \mathcal{M} by restricting it to the elements corresponding to the leaves of a subtree rooted at a vertex v of \mathcal{T} is further denoted by \mathcal{M}_v . The next lemma is the core of the linear-time algorithm for deciding the satisfiability of the formula ψ .

Lemma 8. *Let ψ be a monadic second order logic formula, let \mathcal{T} be a K -decomposition of a matroid \mathcal{M} and let v be a node of \mathcal{T} with children v_1 and v_2 . The ψ -equivalence class of \mathcal{M}_v is uniquely determined by the ψ -equivalence classes of \mathcal{M}_{v_1} and \mathcal{M}_{v_2} , and the functions φ_v and φ_v^r .*

Proof. If the statement of the lemma is false, then there exist \mathcal{M}_{v_1} and \mathcal{M}'_{v_1}

of the same ψ -equivalence class such that the equivalence classes of \mathcal{M}_v and \mathcal{M}'_v are different though \mathcal{M}_{v_2} , φ_v and φ_v^r are the same. Hence, there exists \mathcal{M}_0 such that ψ is satisfied for one of the matroids $\mathcal{M}_v \oplus_K \mathcal{M}_0$ and $\mathcal{M}'_v \oplus_K \mathcal{M}_0$ but not both. Let \mathcal{M} be the matroid such that $\mathcal{M}_{v_1} \oplus_K \mathcal{M} = \mathcal{M}_v \oplus_K \mathcal{M}_0$ and $\mathcal{M}'_{v_1} \oplus_K \mathcal{M} = \mathcal{M}_v \oplus_K \mathcal{M}_0$; such \mathcal{M} exists since it is uniquely determined by \mathcal{M}_0 , \mathcal{M}_{v_2} , φ_v and φ_v^r . By the choice of \mathcal{M}_0 , ψ is satisfied for one of the matroids $\mathcal{M}_{v_1} \oplus_K \mathcal{M}$ and $\mathcal{M}'_{v_1} \oplus_K \mathcal{M}$ but not both and thus \mathcal{M}_{v_1} and \mathcal{M}'_{v_1} cannot be ψ -equivalent as supposed. \square

We are now ready to present the main result of this section.

Theorem 9. *Let ψ be a fixed monadic second order logic and K a fixed integer. There exists an $O(n)$ -time algorithm that given an n -element matroid \mathcal{M} with its K -decomposition decides whether \mathcal{M} satisfies ψ .*

Proof. The core of our algorithm is Lemma 8. Since ψ and K is fixed and the number of ψ -equivalence classes of K -halves is finite by Lemma 7, we can wire in the algorithm the transition table from the equivalence classes of \mathcal{M}_{v_1} and \mathcal{M}_{v_2} , φ_v and φ_v^r to the equivalence class of \mathcal{M}_v where v is a node of \mathcal{T} and v_1 and v_2 its two children. At the beginning, we determine the equivalence classes of \mathcal{M}_v for the leaves v of \mathcal{T} ; this is easy since the equivalence class of \mathcal{M}_v for a leaf v depends only on the fact whether the element corresponding to v is a loop or not.

Then, using the wired in transition table, we determine in constant time the equivalence class of \mathcal{M}_v for each node v based on φ_v , φ_v^r and the equivalence classes of \mathcal{M}_{v_1} and \mathcal{M}_{v_2} for children v_1 and v_2 of v . Once, the equivalence classes of \mathcal{M}_{v_1} and \mathcal{M}_{v_2} for the two children v_1 and v_2 of the root of \mathcal{T} are found, it is easy to determine whether \mathcal{M} satisfies ψ .

As K and ψ are fixed, the running time of our algorithm is clearly linear in the number of nodes of \mathcal{T} which is linear in the number of elements of the matroid \mathcal{M} . \square

Corollary 2 and Theorem 9 yield the following result originally proved by Hliněný [11]:

Corollary 10. *Let \mathbb{F} be a fixed finite field, k a fixed integer and ψ a fixed monadic second order logic formula. There is a polynomial-time algorithm for deciding whether a matroid of branch-width at most k given by its representation over the field \mathbb{F} satisfies ψ .*

7 Deciding representability

We adopt the algorithm presented in [18]. Let us start with recalling some definitions from [18]. A rooted branch-decomposition of \mathcal{M} is obtained from a branch-decomposition of \mathcal{M} by subdividing an edge and rooting the decomposition tree at a new vertex. If \mathcal{M} is a matroid and (A, B) is a partition of its ground set, then two subsets A_1 and A_2 are B -indistinguishable if the following identity holds for every subset B' of B :

$$r(A_1 \cup B') - r(A_1) = r(A_2 \cup B') - r(A_2) .$$

Clearly, the relation of being B -indistinguishable is an equivalence relation on subsets of A . Finally, the transition matrix for an inner node v of a rooted branch-decomposition \mathcal{M} is the matrix whose rows correspond to $\overline{E_1}$ -indistinguishable subsets of E_1 and columns to $\overline{E_2}$ -indistinguishable subsets of E_2 where E_1 and E_2 are the elements corresponding to the leaves of the two subtrees rooted at the children of v and $\overline{E_1}$ and $\overline{E_2}$ are their complements. The entry of \mathcal{M} in the row corresponding to the equivalence class of $F_1 \subseteq E_1$ and in the column corresponding to the equivalence class of $F_2 \subseteq E_2$ is equal to $r(F_1) + r(F_2) - r(F_1 \cup F_2)$. By the definition of indistinguishability, the value of the entry is independent of the choice of F_1 and F_2 in their equivalence classes.

The main algorithmic result of [18] can be reformulated as follows:

Theorem 11. *Let k be a fixed integer and q a fixed prime power. There is a polynomial-time algorithm that for a matroid \mathcal{M} given by its rooted branch-decomposition with transition matrices whose number of rows and columns is at most k and a (oracle-given) mapping of subsets to equivalence classes corresponding to rows and columns of its matrices decides whether \mathcal{M} can be represented over $\mathbb{GF}(q)$ and if so, it computes one of its representations over $\mathbb{GF}(q)$. Moreover, the algorithm can be modified to count all non-isomorphic representations of \mathcal{M} over $\mathbb{GF}(q)$ and to list them (in time linearly dependent on the number of non-isomorphic representations).*

Let \mathcal{T} be a K -decomposition of a matroid \mathcal{M} , v an inner node of \mathcal{T} and E_v the subset of the ground set of \mathcal{M} containing the elements corresponding to the leaves of subtree of \mathcal{T} rooted at v . View \mathcal{T} as a rooted branch-decomposition of \mathcal{M} (we do not claim any upper bound on its width here). Observe that any two subsets of E_v assigned the same color at v are $\overline{E_v}$ -indistinguishable where $\overline{E_v}$ is the complement of E_v . In addition, the values of the function φ_v^r are entries of the transition matrix at v as defined in the beginning of this section. Hence, Theorem 11 yields the following.

Corollary 12. *For every integer K and every prime power q , there is a polynomial-time algorithm that for a matroid \mathcal{M} given by its K -decomposition, decides whether \mathcal{M} can be represented over $\mathbb{GF}(q)$ and if so, it computes one of its representations over $\mathbb{GF}(q)$. Moreover, the algorithm can be modified to count all non-isomorphic representations of \mathcal{M} over $\mathbb{GF}(q)$ and to list them (in time linearly dependent on the number of non-isomorphic representations).*

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