# Facial parity edge colouring

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Abstract: The facial parity edge colouring of a connected bridgeless plane graph is such an edge colouring in which no two face-adjacent edges receive the same colour; in addition, for each face  $\alpha$  and each colour c, no edge or an odd number of edges incident with  $\alpha$  are coloured by c. From the Vizing's theorem it follows that every simple 3-connected plane graph has a such colouring with at most  $\Delta^* + 1$  colours, where  $\Delta^*$  is the size of the largest face. In this paper we prove that any connected bridgeless plane graph has a facial parity edge colouring with at most 92 colours.

Keywords: plane graph, facial walk, edge colouring

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# 1 Introduction

One of the motivations for this paper has come from recent papers of Bunde at al. [1, 2] who introduced parity edge colourings of graphs. Studying the parity of the usage of colours along walks suggested two edge colouring

parameters that have interesting properties and applications. A *parity walk* in an edge colouring of a graph is a walk along which each colour is used an even number of times. Bunde et al. [2] introduced two parameters. Let p(G) be the minimum number of colours in an edge colouring of G having no parity path (a *parity edge colouring*). Let  $\hat{p}(G)$  be the minimum number of colours in an edge colouring of G in which every parity walk is closed (a strong parity edge colouring). Since incident edges of the same colour would form a parity path of length 2, every parity edge colouring is a proper edge colouring, and hence  $p(G) \geq \chi'(G)$ , where  $\chi'(G)$  is the edge chromatic index of G. Since a path is an open walk, any strong parity edge colouring has no parity path. Hence, every strong parity edge colouring is a parity edge colouring and  $\widehat{p}(G) \geq p(G)$  for every graph G. Although there are graphs G with  $\hat{p}(G) > p(G)$  [2], it remains unknown how large  $\hat{p}(G)$  can be when p(G) = k. Elementary results on these parameters appear in [2]. In [1] it is proved that  $\widehat{p}(K_n) = 2^{\lceil \log(n) \rceil} - 1$  for all n. Moreover, the optimal strong parity edge colouring of the complete *n*-vertex graph  $K_n$  is unique when *n* is a power of 2. The authors of [2] mentioned that computing p(G) or  $\widehat{p}(G)$  is NP-hard even when G is a tree. Clearly, the parity edge colouring is such a colouring that each path uses at least one colour an odd number of times.

The vertex version of this problem (strong parity vertex colouring) with some restrictions was introduced in [3]. The authors of [3] conjectured that there is a constant K such that the vertices of any 2-connected plane graph can be coloured with at most K colours in such a way that for each face  $\alpha$ and each colour c, no vertex or an odd number of vertices incident with  $\alpha$  is coloured by c. This conjecture was proved by Kaiser et al. [4], they showed that  $K \leq 300$ , moreover, their colouring is proper.

Another motivation for this work has came from the papers of Pyber [6] and Mátrai [5]. A graph is called *odd* if the degree of its vertices is odd or zero. Pyber raises the problem of edge covering with odd subgraphs in [6] as the counterpart of even subgraph covering problems. He proved that the edges of every finite simple graph can be covered by at most 4 disjoint odd subgraphs; moreover, if the number of vertices is even then 3 colours are sufficient. For not necessarily disjoint coverings we have the following question: Is it true that every graph can be covered by at most 3 odd subgraphs? Mátrai in [5] showed that every finite simple graph can be covered by 3 odd subgraphs and he found an infinite sequence of finite simple connected graphs not coverable by three edge disjoint odd subgraphs.

Pyber's result implies the following: The edges of any 3-connected plane G graph can be coloured by at most 4 colours in such a way that for each face  $\alpha$  and each colour c, no edge or an odd number of edges incident with  $\alpha$  is coloured by c. It is sufficient to consider the dual  $G^*$  of G and its edge

cover with at most 4 disjoint odd subgraphs. This cover induces the required colouring of G.

If we add a requirement that such a colouring must be proper then it is not clear whether there exits such a colouring with K colours, where K is an absolute constant. From Vizing's theorem [7] it follows that every simple 3-connected plane graph G has such a colouring which uses at most  $\Delta^* + 1$ colours, where  $\Delta^*$  is the size of the largest face. Consider the dual graph  $G^*$ and its proper edge colouring. This colouring induces a colouring of G in the natural way. It is such a colouring in which, for each face  $\alpha$  of G, all the edges in the boundary of  $\alpha$  have distinct colours.

From the result of Kaiser et al. [4] follows that each connected bridgeless plane graph has an edge colouring with at most 300 colours such that any two consecutive edges of a facial walk of any face  $\alpha$  receive distinct colours and if a colour *c* appears on a  $\alpha$  then it appears an odd number of times. It is sufficient to consider the medial graph M(G) of *G* and its proper strong parity vertex colouring. The vertices of M(G) correspond to the edges of *G*; two vertices M(G) are adjacent in M(G) if and only if the corresponding edges are face-adjacent in *G*. For any plane *G* the medial graph M(G) is plane as well. Hence, every strong parity vertex colouring of M(G) corresponds to an edge colouring of *G*.

In this paper we show that each connected bridgeless plane graph has such a colouring with at most 92 colours.

# 2 Notation

Let us introduce the notation used in this paper. A graph which can be embedded in the plane is called *planar graph*; a fixed embedding of a planar graph is called *plane graph*.

A *bridge* is an edge whose removal increases the number of components. A graph which contains no bridge is said to be *bridgeless*. In this paper we consider connected bridgeless plane graphs, multiple edges and loops are allowed.

Let G = (V, E, F) be a connected plane graph with the vertex set V, the edge set E, and the face set F. The *degree* of a vertex v, denoted by deg(v), is the number of edges incident with v, each loop counting as two edges. For a face  $\alpha$ , the *size* of  $\alpha$ ,  $deg(\alpha)$ , is defined to be the length of its *facial walk*, i.e. the shortest closed walk containing all edges from the boundary of  $\alpha$ . We often say k-face for a face of size k.

Given a graph G and one of its edge e = uv, (the vertices u and v do not have be different), the *contraction* of e, denoted by G%e, consist of replacing u and v by a new vertex adjacent to all the former neighbours of u and v, and removing the loop corresponding to the edge e. (We keep multiple edges if arisen). Analogously we define the contraction of the set of edges  $H = \{e_1, \ldots, e_k\}$  and we denote it by  $G\%\{e_1, \ldots, e_k\}$  or G%H.

Let H be a subgraph of G. Then the graph  $G \setminus H$  is defined as a graph obtained from G by deleting the vertices in V(H) together their incident edges.

Two faces are *adjacent* if they share an edge. Two (distinct) edges are *face-adjacent* if they are consecutive edges of a facial walk of some face  $\alpha$ .

An edge k-colouring of a graph G is a mapping  $\varphi : E(G) \to \{1, \ldots, k\}$ . The facial parity edge (FPE) colouring of a connected bridgeless plane graph is such an edge colouring that no two face-adjacent edges receive the same colour; for each face  $\alpha$  and each colour c, no edge or an odd number of edges incident with  $\alpha$  are coloured by c.

**Question 1** What is the minimum number of colours  $\chi'_{fp}(G)$  that a connected bridgeless plane graph G has a facial parity edge colouring with at most  $\chi'_{fp}(G)$  colours?

The number  $\chi'_{fp}(G)$  is called the *facial parity chromatic index* of G.

# 3 Results

**Theorem 3.1** Let G be a connected bridgeless plane graph. Then

 $\chi'_{fp}(G) \le 92.$ 

The proof uses the method of discharging. Let G be the counterexample with minimal number of edges, then minimal number of 1-faces, then minimal number of 2-faces. If G is a single cycle of length  $d, d \leq 5$ , we use exactly dcolours. We consider this to be the first step of induction and call this case trivial.

First, we prove several structural properties of G.

We say that a face  $\alpha$  is *small* if  $1 \leq deg(\alpha) \leq 44$  and a face  $\beta$  is *big* if  $deg(\beta) \geq 45$ .

# 3.1 Reducible configurations

We find such (forbidden) subgraphs H of G that the facial parity edge colouring of  $G \setminus H$  or  $G \ H$  using at most 92 colours can be extended to a required colouring of G using at most 92 colours, which is a contradiction to G being a counterexample. In the sequel, whenever we speak about a FPE colouring, we always mean a FPE colouring using at most 92 colours.

#### 3.1.1 1-faces

Claim 1 Each vertex of G is incident with at most one 1-face.

#### $\mathbf{Proof}$

Let v be a vertex incident with at least two 1-faces  $\alpha_1$  and  $\alpha_2$ . If we split v into two vertices  $v_1$  and  $v_2$  in such a way that  $\alpha_1$  and  $\alpha_2$  become a 2-face  $\alpha$ , we obtain a graph G', see Figure 1. It has the same number of edges than G, but less 1-faces. Thus, it is not a counterexample and we can find a FPE colouring  $\varphi'$  of G'. It induces an edge colouring  $\varphi$  of G in a natural way. It is easy to see that  $\varphi$  is a FPE colouring.



Figure 1: A vertex incident with at least two 1-faces can be split into two vertices, reducing the number of 1-faces. Similarly, one can reduce a 1-face and a d-face  $(2 \le d \le 4)$  incident with the same vertex.

**Claim 2** Each vertex of G incident with a 1-face is not incident with any d-face for  $2 \le d \le 4$ .

## Proof

We use the same reduction as in the proof of Claim 1. We split the vertex v incident with a 1-face  $\gamma$  and a d-face  $\alpha$   $(2 \leq d \leq 4)$  in such a way that the faces  $\alpha$  and  $\gamma$  become a (d + 1)-face  $\alpha'$ , see Figure 1 for illustration. Let the reduced graph be G'. It has less 1-faces than G, therefore, it has a FPE colouring  $\varphi'$ . The face  $\alpha'$  has size at most five, therefore, its edges are coloured using d + 1 different colours. Thus, the colouring  $\varphi$  of G induced by the colouring  $\varphi'$  of G' is a FPE colouring, too.

**Claim 3** Each vertex of G incident with a 2-face is not incident with any d-face for  $2 \le d \le 3$ .

#### Proof

We use the same reduction as in the proof of Claims 1 and 2. We omit the details.  $\Box$ 

#### 3.1.2 Small faces

**Claim 4** There are no two small faces adjacent to each other in G.

#### Proof

Let  $\alpha_1$  and  $\alpha_2$  be two small faces adjacent to each other in G.

If both  $\alpha_1$  and  $\alpha_2$  are 1-faces, the graph consists of a single vertex and a loop; it has a FPE colouring using 1 colour. We disregard this trivial case.

Let  $\alpha_1$  be a 1-face and  $\alpha_2$  be a *d*-face,  $d \ge 2$ ; let *e* be the loop they share, see Figure 2. Then the graph G' = G% e has less edges than *G*, therefore, it has a FPE colouring  $\varphi'$ . Let  $\alpha'$  be a face in *G'* corresponding to  $\alpha_1$  and  $\alpha_2$  in *G*. Since  $\alpha_2$  is a small face, at most 43 colours occur on the edges incident with  $\alpha'$ . To extend the colouring  $\varphi'$  of *G'* to a FPE colouring of *G*, it suffices to colour the edge *e* with any colour that does not occur on  $\alpha'$ .

Let  $\alpha_1$  and  $\alpha_2$  be two small faces of size at least 2 and let e be the edge they share, see Figure 2. The graph G' = G% e has less edges than G, therefore, it has a FPE colouring  $\varphi'$ . Let  $\alpha'_1$  and  $\alpha'_2$  be the faces of G' corresponding to the faces  $\alpha_1$  and  $\alpha_2$  in G. (Since  $\alpha_1$  and  $\alpha_2$  are small faces of size at least 2, the size of  $\alpha'_1$  and  $\alpha'_2$  is at least 1.)

Consider the set of colours different from the colours occurring on the edges of  $\alpha'_1$  and  $\alpha'_2$  (the colours admissible for the edge e). Since  $\alpha_1$  and  $\alpha_2$  are small, at most  $2 \cdot (44-1) = 86$  colours occur on their edges. Hence, there is an admissible colour c. We can extend  $\varphi'$  to a FPE colouring  $\varphi$  of G by setting  $\varphi(e) = c$ .



Figure 2: Two adjacent small faces form a reducible configuration: one can contract the edge they share.

**Claim 5** Let  $\beta$  be a big face adjacent to two small faces  $\alpha_1$  and  $\alpha_2$ . Let  $e_i$  be an edge incident with  $\beta$  and  $\alpha_i$ , i = 1, 2. Then  $e_1$  and  $e_2$  are face-adjacent.

## Proof

Let  $e_1$  and  $e_2$  not be face-adjacent. See Figure 3 for illustration. The graph  $G' = G\%\{e_1, e_2\}$  has less edges than G, therefore, it has a FPE colouring  $\varphi'$ . Let  $\alpha'_1, \alpha'_2, \beta'$  be the faces of G' corresponding to the faces  $\alpha_1, \alpha_2, \beta$  in G, respectively.



Figure 3: A big face  $\beta$  adjacent to two different small faces  $\alpha_1$  and  $\alpha_2$  forms a reducible configuration unless the edges  $e_1$  and  $e_2$  are face-adjacent.

We extend the colouring  $\varphi'$  of G' to a FPE colouring of G in the following way: Consider the set of colours different from the colours occurring on the edges of  $\alpha'_1$  and  $\alpha'_2$ ; also different from the colours occurring on the edges of G' corresponding to the edges of G face-adjacent to  $e_1$  and  $e_2$ . There are at least  $92 - 2 \cdot (44 - 1) - 4 = 2$  such colours, say  $c_1$  and  $c_2$ . If at least one of them, say  $c_i$ , already occurs on  $\beta'$ , we set  $\varphi(e_1) = \varphi(e_2) = c_i$ . If none of them occurs on  $\beta'$ , we set  $\varphi(e_i) = c_i$ , i = 1, 2.

Claim 6 Each big face is adjacent to at most one 1-face.

## $\mathbf{Proof}$

It follows from Claims 1 and 5.

Claim 7 Each big face is adjacent to at most two small faces.

#### Proof

Let a big face  $\beta$  be adjacent to small faces  $\alpha_1$ ,  $\alpha_2$ , and  $\alpha_3$ . Consider the edges that  $\beta$  shares with  $\alpha_i$ , i = 1, 2, 3. It is easy to see that there must be a pair of edges  $e_i$  and  $e_j$ , incident to  $\alpha_i$  and  $\alpha_j$ , respectively  $(i \neq j)$ , which are not face-adjacent. It is a contradiction with Claim 5.

#### 3.1.3 Chains of 2-vertices

**Claim 8** There is no chain consisting of at least 5 consecutive 2-vertices in G.

#### Proof

Let  $v_0 e_0 v_1 e_1 \dots v_p e_p v_{p+1}$  be a chain consisting of p vertices of degree 2,  $(v_1, \dots, v_p)$ , where  $p \ge 5$ . The graph  $G' = G\%\{e_1, e_2, e_3, e_4\}$  has a FPE colouring  $\varphi'$ . Let  $e'_0$  and  $e'_5$  be the edges in G' corresponding to the edges  $e_0$  and  $e_5$  in G. Let  $\varphi'(e'_0) = c_1$  and  $\varphi'(e'_5) = c_2$ . The FPE colouring  $\varphi'$  of G' can be extended to a FPE colouring  $\varphi$  of G by setting  $\varphi(e_2) = \varphi(e_4) = c_1$  and  $\varphi(e_1) = \varphi(e_3) = c_2$ , see Figure 4.



Figure 4: A chain of (at least) five 2-vertices is a reducible configuration.

A *d*-face  $\alpha$  is *hanging* on a vertex v, if all vertices incident with  $\alpha$  are 2-vertices except for the vertex v. By Claim 8 we have  $d \leq 5$ . (If deg(v) = 2, the graph G consists of a single cycle of length at most 5, which is the trivial case).

We colour the vertices of G with black, blue and white colour in the following way:

Let all 2-vertices be black, all 3-vertices be blue and all k-vertices for  $k \ge 6$  be white.

A 4-vertex v is black if there is a face hanging on it, else it is white. See Figure 5 for illustration of all types of black vertices.

A 5-vertex v is blue, if there is a face hanging on it, else it is white.

Observe that any black 2-vertex v is incident with two faces. We say v is *bad* for both faces it is incident with. Any black 4-vertex v is incident with a small face  $\alpha$  of size at most 5 and two other faces. The face  $\beta$  adjacent to  $\alpha$  must be big (see Claim 4); the vertex v occurs twice on the facial walk of  $\beta$ . The other face  $\gamma$  can be big or small. We say v is *bad* for the face  $\gamma$ .

Claim 9 For any face  $\alpha$ , there is no chain of at least 5 bad consecutive vertices. Each chain of bad vertices contains at most one bad 4-vertex.



Figure 5: A vertex is black, if it is a 2-vertex, or a 4-vertex with a hanging face.

# Proof

Let  $v_0 e_0 v_1 e_1 v_2 e_2 v_3 e_3 v_4 e_4 v_5 e_5 v_6$  be a subpath of the facial walk of  $\alpha$  containing 5 vertices bad for  $\alpha$ ,  $(v_1, \ldots, v_5)$ . It is easy to see that all the edges  $e_0, \ldots, e_5$  are incident with the same face  $\beta$ . If at least two of the vertices the vertices  $v_1, \ldots, v_5$  are bad 4-vertices, we come to a contradiction with Claim 4 or with Claim 5. If none of them is a 4-vertex, we are in the case of Claim 8. If precisely one of them is a 4-vertex, say  $v_i$ , it is incident with a small face  $\gamma$  of size  $d \leq 5$ .

The graph  $G' = G\%\{e_1, e_2, e_3, e_4\}$  has a FPE colouring  $\varphi'$ . Let  $e'_0$  and  $e'_5$  be the edges in G' corresponding to the edges  $e_0$  and  $e_5$  in G; let  $\alpha', \beta', \gamma'$  be the faces in G' corresponding to the faces  $\alpha, \beta, \gamma$  in G; let v' be the vertex in G' corresponding to  $v_1, \ldots, v_5$  in G. Let  $\varphi'(e'_0) = c_1$  and  $\varphi'(e'_5) = c_2$ . Since  $e'_0$  and  $e'_5$  are face-adjacent in G',  $c_1 \neq c_2$ . To extend  $\varphi'$  to a FPE colouring of G, we proceed in the following way:

If i = 1, 3, or 5, we simply set  $\varphi(e_2) = \varphi(e_4) = c_1$  and  $\varphi(e_1) = \varphi(e_3) = c_2$ .

If i = 2 or i = 4, we set  $\varphi(e_2) = \varphi(e_4) = c_1$  and  $\varphi(e_1) = \varphi(e_3) = c_2$ and switch the order of colours of edges incident with  $\gamma'$ , see Figure 6 for illustration (here the small face  $\gamma$  is a bigon.).

**Claim 10** Let v be a black vertex bad for a small face  $\alpha$ . If v is a 4-vertex, then the face hanging on it is a 1-face.

# Proof

Let v be a black 4-vertex, let  $\gamma$  be the face hanging on v of size at least 2 and let  $e_1$  and  $e_2$  be the edges of  $\alpha$  incident with v. It is easy to see that the edges  $e_1$  and  $e_2$  are incident with the same big face, say  $\beta$ .

There is an edge  $e_{\gamma}$  incident with  $\gamma$  and  $\beta$ , which is not face-adjacent to  $e_{\alpha} \in \{e_1, e_2\}$ , what is a contradiction with Claim 5.

**Claim 11** Let  $\alpha$  be a small face sharing at least two bad black vertices with a big face  $\beta$ . Then all the bad black vertices incident with  $\alpha$  and  $\beta$  are 2-vertices.



Figure 6: A chain of (at least) five bad black vertices is a reducible configuration as well.

# Proof

Let  $v_1$  and  $v_2$  be black vertices incident with  $\alpha$  and  $\beta$ . If  $v_1$  is a 4-vertex, it is bad for  $\alpha$ , therefore, there is a small face  $\gamma$  hanging on  $v_1$ , adjacent to  $\beta$ . By Claim 10, the face  $\gamma$  is a 1-face. Thus, we can find an edge (a loop)  $e_{\gamma}$ incident with  $\gamma$  and  $\beta$  and an edge  $e_{\alpha}$  incident with  $\alpha$  and  $\beta$ , which are not face-adjacent. It is a contradiction with Claim 5.

Claim 12 Let v be a black vertex bad for a small d-face  $\alpha$ ,  $d \in \{2, 3, 4\}$ . Then v is a 2-vertex.

## Proof

Let v be a black 4-vertex, let  $\gamma$  be the face hanging on v. By Claim 10, the face  $\gamma$  is a 1-face. On the other hand, by Claim 2, the face  $\gamma$  cannot be a 1-face, what is a contradiction.

**Claim 13** If a big face  $\beta$  shares a 2-vertex v with a small face  $\alpha$ , then  $\beta$  is not adjacent to any other small face.



Figure 7: Reducible pairs of edges  $(e_{\gamma} \text{ and } e_{\alpha})$ . For details see Claims 10 and 11.

## Proof

It follows immediately from Claim 5.

**Claim 14** Let  $\gamma$  be a d-face,  $d \in \{2, 3, 4, 5\}$ , hanging on a vertex v, adjacent to a big face  $\beta$ . Then  $\beta$  is not adjacent to any other small face.

## Proof

It follows immediately from Claim 13.

# 3.2 Discharging rules

Let G be a counterexample. It contains no reducible configuration. Let the initial charge of each vertex be  $\psi(v) = \deg(v) - 6$  and the initial charge of each face be  $\psi(\alpha) = 2 \deg(\alpha) - 6$ . From Euler's formula we can easily derive that

$$\sum_{\alpha \in F} (2\deg(\alpha) - 6) + \sum_{v \in V} (\deg(v) - 6) = -12.$$

It is obvious that all the negative charge is in the vertices of degree 2, 3, 4, and 5 and in the faces of size 1 and 2.

**Rule 1:** Let  $\beta$  be a big face.

- If  $\beta$  is adjacent to a single small face  $\alpha$ , it sends 3 units of charge to  $\alpha$ .
- If  $\beta$  is adjacent to two small faces  $\alpha_1$  and  $\alpha_2$ , such that  $\deg(\alpha_1) \leq \deg(\alpha_2)$ , it sends 2 units of charge to  $\alpha_1$  and 1 unit of charge to  $\alpha_2$ . (If  $\deg(\alpha_1) = \deg(\alpha_2)$ , we do it arbitrarily.)

**Rule 2:** Let  $\beta$  be a big face.

- It sends 2 units of charge to any black vertex bad for  $\beta$ .
- It sends 1 unit of charge to any other black, blue or white vertex incident with β. (Multiply incident vertices are considered as different).

**Rule 3:** Let  $\alpha$  be a small face.

- It sends 2 units of charge to any black vertex bad for  $\alpha$ .
- It sends 1 unit of charge to any other black or blue vertex incident with  $\alpha$ .

**Rule 4:** Let v be a black 4-vertex.

• It sends 2 units of charge to the incident small hanging face  $\gamma$ .

**Rule 5:** Let v be a blue 5-vertex.

• It sends 2 units of charge to the incident small hanging face  $\gamma$ .

**Rule 6:** Let v be a k-vertex,  $k \ge 6$ .

• It sends 2 units of charge to any incident small hanging face  $\gamma$ .

# 3.3 Analysis of the graph

#### 3.3.1 Vertices

Every 2-vertex is bad black for both faces incident with it, hence it receives 2 units of charge from both incident faces (Rules 2 and 3). Its charge is -4 + 2 + 2 = 0.

Every 3-vertex is blue, hence it receives 1 unit of charge from all the three incident faces (Rules 2 and 3). Its charge is  $-3 + 3 \cdot 1 = 0$ .

Every black 4-vertex v receives 2 units of charge from the face  $\alpha$  it is bad for (Rules 2 and 3) and  $2 \cdot 1$  units of charge from the doubly-incident big face  $\beta$  (Rule 2). It sends 2 units of charge to the hanging face  $\gamma$  (Rule 4). The charge of v is  $-2 + 2 + 2 \cdot 1 - 2 = 0$ .

Every white 4-vertex is incident with at least 2 big faces (see Claim 4), therefore, its charge is at least  $-2 + 2 \cdot 1 = 0$ .

Every blue 5-vertex v is incident with a hanging face  $\gamma$ , doubly-incident with a big face  $\beta$  and incident with two more faces  $\alpha_1$  and  $\alpha_2$ . Therefore, v receives  $4 \cdot 1$  units of charge from the incident faces (Rules 2 and 3). It sends 2 units of charge to  $\gamma$  (Rule 5). Therefore, the charge of v is  $-1+4\cdot 1-2=1$ .

Every white 5-vertex is incident with at least 3 big faces (see Claim 4), therefore, its charge is at least  $-1 + 3 \cdot 1 = 2$ .

Every (white) k-vertex  $v, k \ge 6$ , has non-negative initial charge. It receives charge from big faces (Rule 2) and sends charge to the hanging faces  $\gamma_1, \ldots, \gamma_r$  (Rule 6). For each hanging face  $\gamma_i$ , the adjacent big face  $\beta_i$  is doubly-incident to v. The faces  $\beta_i$  and  $\beta_j$  are different for different  $\gamma_i$  and  $\gamma_j$ (see Claims 5 and 6). Therefore, the charge of v is at least  $r \cdot (2 \cdot 1 - 2) = 0$ .

## 3.3.2 1-faces

Let  $\gamma$  be a 1-face. It is adjacent to a big face  $\beta$  and it is hanging on a vertex v. The face  $\beta$  is adjacent to at most one 1-face (Claim 6). Therefore, it sends at least 2 units of charge to the face  $\gamma$  (Rule 1).

If the vertex v is a 2-vertex, we get the trivial graph, which was omitted before. If the vertex v is a 3-vertex, the third edge incident with v is a bridge in G, which is not allowed. Therefore,  $\deg(v) \ge 4$  and the vertex v sends 2 units of charge to the face  $\gamma$  (Rules 4 - 6). The charge of  $\gamma$  is at least -4 + 2 + 2 = 0.

# 3.3.3 2-faces

Let  $\alpha$  be a 2-face. Its initial charge is -2. Let  $v_1$  and  $v_2$  be the vertices incident with  $\alpha$ . The face  $\alpha$  is adjacent to at most 2 faces, which must be big (see Claim 4). Consider the number of black vertices bad for  $\alpha$ . Note that by Claim 12 each such vertex is a 2-vertex.

- 1. Let both  $v_1$  and  $v_2$  be bad black. Then the graph G consists of a single cycle on 2 vertices, which is not a counterexample.
- 2. Let  $v_1$  be a black 2-vertex. Then  $\alpha$  is adjacent to a single big face  $\beta$ . The big face  $\beta$  is not adjacent to any other small face (see Claim 13). Therefore,  $\beta$  sends 3 units of charge to  $\alpha$  (Rule 1). The face  $\alpha$  sends 2 units of charge to  $v_1$  and at most 1 unit of charge to  $v_2$  (Rule 3). On the other hand,  $\alpha$  is hanging on  $v_2$ , therefore, it receives 2 units of charge from  $v_2$  (Rules 4 - 6). Note that  $v_2$  cannot be a 3-vertex, otherwise there would be a bridge in G. The charge of  $\alpha$  is at least -2 - 2 - 1 + 2 + 3 = 0.
- 3. Let none of  $v_1$  and  $v_2$  be bad black. Consider the number of faces adjacent to  $\alpha$ . If  $\alpha$  is adjacent to a single big face  $\beta$ , then  $\beta$  is not adjacent to any other small face. Therefore,  $\beta$  sends 3 units of charge to  $\alpha$ . Moreover, in this case none of  $v_1$  and  $v_2$  can be black nor blue. The charge of  $\alpha$  is -2 + 3 = 1.

If  $\alpha$  is adjacent to two big faces  $\beta_1$  and  $\beta_2$ , the face  $\alpha$  sends at most  $2 \cdot 1$  unit of charge to  $v_1$  and  $v_2$ . The big face  $\beta_i$ , i = 1, 2, is not adjacent to any other small face of size at most 2 (see Claims 5, 2, and 3). Therefore, by Rule 1, the face  $\beta_i$ , i = 1, 2, sends at least 2 units of charge to  $\alpha$ . The charge of  $\alpha$  is at least  $-2 - 2 \cdot 1 + 2 \cdot 2 = 0$ .



Figure 8: Different possible neighbourhoods of a 2-face  $\alpha$ .

# 3.3.4 3-faces

Let  $\alpha$  be a 3-face. Its initial charge is 0. Let  $v_1, v_2, v_3$  be the vertices incident with  $\alpha$ . (They do not have to be pairwise different.) Consider the number of black vertices bad for  $\alpha$ . Note that by Claim 12 each such vertex is a 2-vertex.

- 1. Let all the three vertices  $v_1, v_2, v_3$  be bad black. Then the graph G consists of a single cycle on 3 vertices, which is not a counterexample.
- 2. Let  $v_2$  and  $v_3$  be bad black. Then all the three edges of  $\alpha$  are incident with the same big face  $\beta$  and  $\alpha$  is hanging on  $v_1$ . Then by Claim 14 the face  $\beta$  sends 3 units of charge to  $\alpha$  (Rule 1). The face  $\alpha$  then sends  $2 \cdot 2$  units of charge to the 2-vertices  $v_2, v_3$  and at most 1 unit of charge to the vertex  $v_1$  (Rule 3). If the vertex  $v_1$  is a 3-vertex, the third edge incident with it is a bridge in G. Therefore deg $(v_1) \geq 4$  and  $v_1$  sends 2 units of charge to  $\alpha$  (Rules 4-6). The charge of  $\alpha$  is  $3-2\cdot 2-1+2=0$ .
- 3. Let  $v_3$  be bad black. Let  $\beta_1$  be the big face incident with the edge  $v_1v_2$ and  $\beta_2$  be the big face incident with the edges  $v_2v_3$  and  $v_3v_1$ .

The face  $\beta_2$  sends 3 units of charge to  $\alpha$ ,  $\beta_1$  sends at least 1 unit of charge to  $\alpha$ . The face  $\alpha$  then sends 2 units of charge to  $v_3$  and at most 1 unit of charge to  $v_1$  and  $v_2$ . The charge of  $\alpha$  is at least  $3+1-2-2\cdot 1=0$ .

4. Let none of the vertices  $v_1, v_2, v_3$  be bad black. Consider the number of faces adjacent to  $\alpha$ .

If there are three different faces adjacent to  $\alpha$ , (they must be big, see Claim 4), the face  $\alpha$  receives at least 1 unit of charge from each of them, and sends at most 1 unit of charge to each incident vertex. Hence, the charge of  $\alpha$  is at least 0.

If there is a big face  $\beta$  sharing at least two edges with  $\alpha$ , these edges are not face-adjacent in  $\beta$ . Therefore,  $\beta$  is not adjacent to any other

small face, thus, it sends 3 units of charge to  $\alpha$ . The charge of  $\alpha$  is non-negative again.



Figure 9: Different possible neighbourhoods of a 3-face  $\alpha$ .

# 3.3.5 4-faces

Let  $\alpha$  be a 4-face. Its initial charge is 2. Let  $v_1, v_2, v_3, v_4$  be the vertices incident with  $\alpha$ . (They do not have to be pairwise different.) Consider the number of black vertices bad for  $\alpha$ . Note that by Claim 12 each such vertex is a 2-vertex.

- 1. Let all the four vertices  $v_1, v_2, v_3, v_4$  be bad black. Then the graph G consists of a single cycle on 4 vertices, which is not a counterexample.
- 2. Let  $v_1, v_2$ , and  $v_3$  be bad black. Then all the four edges of  $\alpha$  are incident with the same big face  $\beta$  and  $\alpha$  is hanging on  $v_4$ . Hence,  $\beta$  sends 3 units of charge to  $\alpha$  (Rule 1). The face  $\alpha$  then sends  $3 \cdot 2$  units of charge to the 2-vertices  $v_1, v_2, v_3$  and at most one unit of charge to  $v_4$  (Rule 3). If the vertex  $v_4$  is a 3-vertex, the third edge incident with it is a bridge in G. Therefore, deg $(v_4) \geq 4$  and  $v_4$  sends 2 units of charge to  $\alpha$  (Rules 4-6). The charge of  $\alpha$  is  $2+3-3\cdot 2-1+2=0$ .
- 3. Let  $v_1$  and  $v_3$  be bad black. Let  $\beta_1$  be the big face incident with  $v_1$ , let  $\beta_2$  be the big face incident with  $v_3$ .

If  $\beta_1 \neq \beta_2$ , both  $\beta_1$  and  $\beta_2$  send 3 units of charge to  $\alpha$ . The face  $\alpha$  then sends 2 units of charge to the vertices  $v_1$  and  $v_3$  and at most 1 unit of charge to the vertices  $v_2$  and  $v_4$ . The charge of  $\alpha$  is at least  $2 + 2 \cdot 3 - 2 \cdot 2 - 2 \cdot 1 = 2$ . If  $\beta_1 = \beta_2$ , then  $v_2$  and  $v_4$  are not blue, and the charge of  $\alpha$  is at least  $2 + 3 - 2 \cdot 2 = 1$ .

4. Let  $v_1$  and  $v_2$  be bad black. Let  $\beta_1$  be the big face incident with the edge  $v_3v_4$  and  $\beta_2$  be the big face incident with the vertices  $v_1$  and  $v_2$ . If  $\beta_1 \neq \beta_2$ ,  $\beta_2$  sends 3 units of charge to  $\alpha$  and  $\beta_1$  sends at least 1 unit of

charge to  $\alpha$ . The face  $\alpha$  then sends 2 units of charge both to  $v_1$  and  $v_2$ and at most 1 unit of charge to  $v_3$  and  $v_4$ . The charge of  $\alpha$  is at least  $2+3+1-2\cdot 2-2\cdot 1=0$ . If  $\beta_1=\beta_2$ , then  $v_3$  and  $v_4$  are not blue, and the charge of  $\alpha$  is at least  $2+3-2\cdot 2=1$ .

- 5. Let  $v_1$  be bad black. The big face  $\beta_1$  incident with  $v_1$  sends 3 units of charge to  $\alpha$ . The face  $\alpha$  then sends 2 units of charge to  $v_1$  and at most 1 unit of charge to  $v_2, v_3$ , and  $v_4$ . The charge of  $\alpha$  is at least  $2+3-2-3\cdot 1=0$ .
- 6. Let no bad black vertex be incident with  $\alpha$ . Then the big faces adjacent to  $\alpha$  send together at least 2 units of charge (if there was only one big face, it would send 3 units of charge), and  $\alpha$  sends at most 1 unit of charge to each incident vertex. The charge of  $\alpha$  is at least 2+2-4=0.



Figure 10: Different possible neighbourhoods of a 4-face  $\alpha$ .

#### 3.3.6 5-faces

Let  $\alpha$  be a 5-face. Its initial charge is 4. Let  $v_1, v_2, v_3, v_4, v_5$  be the vertices incident with  $\alpha$ . (They do not have to be pairwise different.) Consider the number of black vertices bad for  $\alpha$ . Note that by Claim 10 each such vertex is either a 2-vertex or a 4-vertex with a hanging 1-face.

- 1. Let all the five vertices  $v_1, v_2, v_3, v_4, v_5$  be bad black. Then the graph G contains only 5 vertices, hence, it is not a counterexample.
- 2. Let  $v_1, v_2, v_3$ , and  $v_4$  be bad black. From Claim 5 follows that none of them is incident with a 1-face. Then all the five edges of  $\alpha$  are incident with the same big face  $\beta$  and  $\alpha$  is hanging on  $v_5$ . Hence,  $\beta$  sends 3 units of charge to  $\alpha$  (Rule 1). The face  $\alpha$  then sends  $4 \cdot 2$  units of charge to the 2-vertices  $v_1, v_2, v_3$ , and  $v_4$  (Rule 3) and at most 1 unit of charge to the vertex  $v_5$  (Rule 3). If the vertex  $v_5$  is a 3-vertex, the third edge incident with it is a bridge in G. Therefore deg $(v_5) \geq 4$

and  $v_5$  sends 2 units of charge to  $\alpha$  (Rules 4 – 6). The charge of  $\alpha$  is  $4+3-4\cdot 2-1+2=0$ .

- 3. Let  $v_1, v_2$ , and  $v_3$  be bad black. Let  $\beta_1$  be the big face incident with the vertices  $v_1, v_2$ , and  $v_3$  and  $\beta_2$  be the big face incident with the edge  $v_4v_5$ . By Claim 11  $v_1, v_2$ , and  $v_3$  are 2-vertices. If  $\beta_1 \neq \beta_2$ ,  $\beta_1$  sends 3 units of charge to  $\alpha$  and  $\beta_2$  sends at least 1 unit of charge to  $\alpha$ . The face  $\alpha$  then sends 2 units of charge to  $v_1, v_2$ , and  $v_3$  and at most 1 unit of charge to  $v_4$  and  $v_5$ . The charge of  $\alpha$  is at least  $4 + 3 + 1 - 3 \cdot 2 - 2 \cdot 1 = 0$ . If  $\beta_1 = \beta_2$ , then  $v_4$  and  $v_5$  are not blue, and the charge of  $\alpha$  is at least  $4 + 3 - 3 \cdot 2 = 1$ .
- 4. Let  $v_1, v_2$ , and  $v_4$  be bad black. Let  $\beta_1$  be the big face incident with  $v_1$ and  $v_2$ , let  $\beta_2$  be the big face incident with  $v_4$ . By Claim 11  $v_1$  and  $v_2$ are 2-vertices, thus  $\beta_1$  sends 3 units of charge to  $\alpha$ .

If  $\beta_1 \neq \beta_2$ , the face  $\alpha$  sends 2 units of charge to the vertices  $v_1, v_2$ , and  $v_4$  and at most 1 unit of charge to the vertices  $v_3$  and  $v_5$ . The charge of  $\alpha$  is at least  $4 + 3 + 1 - 3 \cdot 2 - 2 \cdot 1 = 0$ . If  $\beta_1 = \beta_2$ , then  $v_3$  and  $v_5$  are not blue, and the charge of  $\alpha$  is  $4 + 3 - 3 \cdot 2 = 1$ .

- 5. Let  $v_1$  and  $v_2$  be bad black. Let  $\beta_1$  be the big face incident with the vertices  $v_1$  and  $v_2$ . The face  $\beta_1$  sends 3 units of charge to  $\alpha$ . The face  $\alpha$  then sends 2 units of charge both to  $v_1$  and  $v_2$  and at most 1 unit of charge to  $v_3, v_4$ , and  $v_5$ . The charge of  $\alpha$  is at least  $4+3-2\cdot 2-3\cdot 1=0$ .
- 6. Let  $v_1$  and  $v_3$  be bad black. Let  $\beta_1$  be the big face incident with  $v_1$ ,  $\beta_2$  be the big face incident with  $v_3$ , and  $\beta_3$  be the big face incident with the edge  $v_4v_5$ . If  $\beta_1, \beta_2$ , and  $\beta_3$  are three different faces, the charge of  $\alpha$  is at least  $4 + 3 2 \cdot 2 3 \cdot 1 = 0$ . If two of them coincide, then at least one of the vertices  $v_2, v_4$ , and  $v_5$  is not blue and the charge of  $\alpha$  is at least  $4 + 2 2 \cdot 2 2 \cdot 1 = 0$ . If  $\beta_1 = \beta_2 = \beta_3$ , then  $v_2, v_4$ , and  $v_5$  are not blue and the charge of  $\alpha$  is at least  $4 + 2 2 \cdot 2 2 \cdot 1 = 0$ . If  $\beta_1 = \beta_2 = \beta_3$ , then  $v_2, v_4$ , and  $v_5$  are not blue and the charge of  $\alpha$  is at least  $4 + 1 2 \cdot 2 = 1$ .
- 7. Let  $v_1$  be black. Then the big faces adjacent to  $\alpha$  send together at least 2 units of charge (if there was only one big face, it would send 3 units of charge). The charge of  $\alpha$  is at least  $4 + 2 2 4 \cdot 1 = 0$ .
- 8. Let no black vertex be incident with  $\alpha$ . Then the big faces adjacent to  $\alpha$  send together at least 1 unit of charge. The charge of  $\alpha$  is at least 4+1-5=0.



Figure 11: Different possible neighbourhoods of a 5-face  $\alpha$ .

#### 3.3.7 Small faces of size at least 6

Let  $\alpha$  be a *d*-face,  $6 \leq d \leq 44$ . Its initial charge is 2d - 6. Let  $v_1, \ldots, v_d$  be the vertices incident with  $\alpha$ . (They do not have to be pairwise different.) Consider the black vertices incident with  $\alpha$ . Let  $v_i$  be a black 4-vertex. It cannot be good for  $\alpha$ , since no two small faces are adjacent (see Claim 4). Therefore, each black 4-vertex is bad for  $\alpha$ . By Claim 9 at most d - 2 vertices incident with  $\alpha$  are bad. Let  $k \leq d - 2$  be the number of black vertices incident with  $\alpha$ . We can divide the facial walk of  $\alpha$  into  $d - k \geq 2$  parts, each beginning and ending in a blue or white vertex, each incident with  $\alpha$  and a big face  $\beta_i, i \in \{1, \ldots, d - k\}$ . Each of these big faces sends at least 1 unit of charge to  $\alpha$ . (If  $\beta_i = \beta_j$  for some  $1 \leq i < j \leq d - k$ , the face  $\beta_i$  cannot be adjacent to other small face than  $\alpha$ , therefore, it sends 3 units to  $\alpha$ , which is even more than what two different big faces would send.)

The face  $\alpha$  then sends 2 units of charge to each of the k incident black vertices, and at most 1 unit of charge to each of the other incident vertices. Together, the charge of  $\alpha$  is at least

$$2d - 6 + (d - k) \cdot 1 - k \cdot 2 - (d - k) \cdot 1 = 2(d - k) - 6.$$

If  $d - k \ge 3$ , the charge of  $\alpha$  is non-negative.

Let d - k = 2. It means there are only two vertices which are not black. Since  $d \ge 6$ , at least one big face  $\beta$  shares at least 2 black vertices with  $\alpha$ , say  $v_1$  and  $v_2$ . By Claim 9 at least one from  $v_1$  and  $v_2$  is a 2-vertex, hence, by Claim 13 the face  $\beta$  sends 3 units of charge to  $\alpha$ . The charge of  $\alpha$  is therefore at least

$$2d - 6 + 3 + 1 - (d - 2) \cdot 2 - 2 \cdot 1 = 0.$$

#### 3.3.8 Big faces

Let  $\beta$  be a *d*-face,  $d \geq 45$ . Its initial charge is 2d - 6. It sends 3 units of charge to the small faces it is adjacent to (Rule 1). It sends 2 units of charge to all bad black vertices; 1 unit of charge to all other vertices. Let k be the number of black vertices bad for  $\beta$ . By Claim 9,  $k \leq \frac{4}{5} \cdot d$ . The charge of  $\beta$  is therefore at least

$$2d - 6 - 3 - k \cdot 2 - (d - k) \cdot 1 = d - k - 9 \ge d - \frac{4d}{5} - 9 = \frac{d}{5} - 9 = \frac{d - 45}{5} \ge 0.$$

The charge of all elements of the graph is non-negative, but the sum of all the charge is -12, what is a contradiction.

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