Gray Codes Faulting Matchings

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Abstract

A (cyclic) n-bit Gray code is a (cyclic) ordering of all 2^n binary strings of length n such that consecutive strings differ in a single bit. Equivalently, an n-bit Gray code can be viewed as a Hamiltonian path of the n-dimensional hypercube Q_n, and a cyclic Gray code as a Hamiltonian cycle of Q_n. In this paper we study Hamiltonian paths and cycles of Q_n avoiding a given set of faulty edges that form a matching, briefly called (cyclic) Gray codes faulting a given matching. Given a matching M and two vertices u, v of Q_n, $n \geq 4$, our main result provides a necessary and sufficient condition, expressed in terms of forbidden configurations for M, for the existence of a Gray code between u and v faulting M. As a corollary, we obtain a similar characterization for a cyclic Gray code faulting M. In particular, in case that M is a perfect matching, Q_n has a (cyclic) Gray code faulting M if and only if $Q_n - M$ is a connected graph. This complements a recent result of Fink, who proved that every perfect matching of Q_n can be extended to a Hamiltonian cycle.

Keywords: hypercube, Gray code, Hamiltonian paths, Hamiltonian cycles, matching

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1 Introduction

A (*cyclic*) *n-bit Gray code* is a (cyclic) ordering of all 2^n binary strings of length n such that consecutive strings differ in a single bit. It is named after Frank Gray, who in 1953 patented a simple scheme to generate such a cyclic code for every $n \geq 2$ [7]. Since then, the research on Gray codes satisfying certain additional properties has received a considerable attention, and applications have been found in such diverse areas as data compression, graphics and image processing, information retrieval, signal encoding or processor allocation in hypercubic networks [12].

Alternatively, an n-bit Gray code can be viewed as a Hamiltonian path of the n-dimensional hypercube Q_n, and a cyclic Gray code as a Hamiltonian cycle of Q_n. The applications of hypercubes in parallel computing [10] inspired the investigation of Hamiltonian paths and cycles in hypercubes avoiding a given set of faulty edges. Chan and Lee [2] proved that the problem whether Q_n contains a Hamiltonian cycle avoiding a given set $\mathcal{F}$ of faulty edges is NP-complete. On the other hand, they showed that if $|\mathcal{F}| \leq 2n - 5$, $n \geq 3$, then such a Hamiltonian cycle exists if and only if each vertex is incident with at least two nonfaulty edges. This upper bound is sharp in the sense that for any $n \geq 3$ there are $2n-4$ faulty edges in Q_n satisfying the above condition, but there is no Hamiltonian cycle avoiding them. Tsai [13] obtained a similar result for Hamiltonian paths with given endvertices.

A related problem considers Hamiltonian cycles and paths in hypercubes passing through a set of prescribed edges. Caha and Koubek [1] observed that if a set $\mathcal{P}$ of prescribed edges extends to a Hamiltonian cycle, then it induces a subgraph consisting of pairwise vertex-disjoint paths. On the other hand, they showed that for any $n \geq 3$ there is a set of $2n - 2$ edges, satisfying this necessary condition, but there is no Hamiltonian cycle of Q_n containing them. This bound is sharp in the sense that if $|\mathcal{P}| \leq 2n - 3$, $n \geq 2$, the necessary condition is also sufficient for a Hamiltonian cycle of Q_n passing through every edge of $\mathcal{P}$ [3]. There is a similar result for Hamiltonian paths with prescribed endvertices if $|\mathcal{P}| \leq 2n - 4$, $n \geq 5$ [4]. A special case of the above problem is when the prescribed edges comprise a matching of the hypercube. Kreweras [9] conjectured that every perfect matching of Q_n, $n \geq 2$, can be completed to a Hamiltonian cycle. Recently, the conjecture was proved by Fink [5].

In this paper we study Hamiltonian paths and cycles in Q_n avoiding a given set of faulty edges that form a matching, or equivalently, (cyclic) Gray codes faulting a given matching. Given a matching M and two vertices u, v of Q_n, $n \geq 4$, our main result provides a necessary and sufficient condition for the existence of a Gray code between u and v that faults M. As a corollary, we obtain a similar characterization for the existence of a cyclic Gray code faulting M. In particular, in case that M is a perfect matching, Q_n has a (cyclic) Gray code faulting M if and only if $Q_n - M$ is a connected graph.

Note that this complements the above quoted result of Fink. Furthermore, since our necessary and sufficient conditions are decidable in a polynomial time, our results imply that the problem of Hamiltonicity of Q_n with faulty edges, which is NP-complete in general, becomes polynomial for up to 2^{n-1} edges provided they form a matching.

After definitions and preliminary results in the next two sections, we study necessary conditions in Sections 4 and 5. Then we consider the initial case of $n = 3$ and $n = 4$ in Sections 6 and 7. The induction step is presented in Section 8. The last section summarizes obtained results.

2 Preliminaries

In this text n is always a positive integer and $[n]$ denotes the set $\{1, 2, \ldots, n\}$. As usual, the vertex and edge sets of a graph G are denoted by $V(G)$ and $E(G)$, respectively. For a vertex $v \in V(G)$ and a set $E \subseteq E(G)$, let $G - v$ denote the subgraph of G induced by $V(G) \setminus \{v\}$, and let $G - E$ denote the graph with vertices $V(G)$ and edges $E(G) \setminus E$. Similarly, for a pair of vertices $x, y \in V(G)$ let $G + xy$ denote the graph with vertices $V(G)$ and edges $E(G) \cup \{xy\}$. The distance of vertices u and v is denoted by $d(u, v)$. The distance $d(u, vw)$ of a vertex u and an edge vw is defined as the minimum of distances $d(u, v)$ and $d(u, w)$.

A path with endvertices a and b is denoted by P_{ab}. In particular, P_{aa} denotes the path consisting of a single vertex a. If P_{ab} and P_{cd} are vertex disjoint paths and b and c are adjacent vertices, then $P_{ab} + P_{cd}$ denotes the path with endvertices a and d, obtained as a concatenation of P_{ab} with P_{cd}. A path P is called a *subpath* of a path P' if P forms a subgraph of P'. We say that paths $\{P_i\}_{i=1}^n$ are *spanning paths* of a graph G if $\{V(P_i)\}_{i=1}^n$ partitions $V(G)$.

The *n-dimensional hypercube* Q_n is the graph with vertex set $V(Q_n) = \{0, 1\}^n$ and edge set $E(Q_n) = \{uv \mid |\Delta(u, v)| = 1\}$ where $\Delta(u, v) = \{i \in [n] \mid u_i \neq v_i\}$. The *dimension* $\dim(uv)$ of an edge $uv \in E(Q_n)$ is the integer d such that $\Delta(u, v) = \{d\}$. For a vertex $v \in V(Q_n)$ let v^d denote the vertex of Q_n such that vv^d is an edge of dimension d.

All edges of the same dimension d form a *layer* (*of dimension* d) in Q_n. Note that $E(Q_n)$ is partitioned into n layers, each of size 2^{n-1}. Moreover, for every $d \in \Delta(u, v)$, each path P_{uv} contains an edge of dimension d . In other words, vertices u and v are separated by the layer of dimension d.

The *parity* $p(v)$ of a vertex $v \in V(Q_n)$ is defined by $p(v) = \|v\| \bmod 2$ where $\|v\| = \sum_{i=1}^{n} v_i$. Note that vertices of each parity form bipartite sets of Q_n. Consequently, $p(u) = p(v)$ if and only if $d(u,v)$ is even. We extend the definition of the parity to edges by putting for $uv \in E(Q_n)$

$$p(uv) = \begin{cases} p(u), & \text{if } \|u\| < \|v\|; \\ p(v), & \text{otherwise.} \end{cases}$$

For every $d \in [n]$ let Q_L^d and Q_R^d denote the subgraphs of Q_n induced by the sets $\{v \in V(Q_n) \mid v_d = 0\}$ and $\{v \in V(Q_n) \mid v_d = 1\}$, respectively. Note that both Q_L^d and Q_R^d are isomorphic to Q_{n-1}. Similarly, for a matching M in Q_n put

$$\begin{aligned} M^d &= \{e \in M \mid \dim(e) = d\}, \\ M_L^d &= M \cap E(Q_L^d), \text{ and} \\ M_R^d &= M \cap E(Q_R^d). \end{aligned}$$

Observe that $(Q_i^a)_j^b = (Q_j^b)_i^a$ and $(M_i^a)_j^b = (M_j^b)_i^a$ for every $a, b \in [n]$ and $i, j \in \{L, R\}$.

Given a set F of faulty edges of Q_n, we say that a subgraph of Q_n *faults* F if it forms a subgraph of $Q_n - F$.

3 Some fundamental results

A classical result of Lewinter and Widulski [11] describes Hamiltonian paths in hypercubes with one faulty vertex.

Lemma 3.1. *For any pairwise distinct vertices u, v, w in Q_n such that $p(u) = p(v) \neq p(w)$, there exists a Hamiltonian path of $Q_n - w$ between u and v.*

There is a similar characterization for hypercubes with at most one faulty edge, which follows from more general results of Tsai et al. [14]. It should be remarked that the special case of $M = \emptyset$ was firstly observed by Havel [8].

Lemma 3.2. *For every $n \geq 3$, vertices u and v of different parities, and a matching M of Q_n with $|M| \leq 1$, there exists a Hamiltonian path in $Q_n - M$ between u and v.*

An extension of the above mentioned Havel's result was obtained in [3].

Lemma 3.3. *For every distinct vertices u, v, x, y of Q_n such that $p(u) \neq p(v)$ and $p(x) \neq p(y)$, there exist spanning paths P_{uv}, P_{xy} of Q_n.*

We shall also employ the following lemma on the existence of a Hamiltonian path passing through a prescribed edge, which is a special case of more general results of Caha and Koubek [1].

Lemma 3.4. *For every two vertices u, v and an edge e of Q_n such that $p(u) \neq p(v)$ and $e \neq uv$, there exists a Hamiltonian path of Q_n between u and v passing through e.*

We conclude this section with a simple observation.

Proposition 3.5. *For every matching M in Q_n with $|M| = 2$ there exists a dimension $d \in [n]$ such that $|M^d_L| = |M^d_R| = 1$.*

Proof. We argue by induction on n. Since the case $n \leq 2$ is obvious, assume that $n \geq 3$ and choose an arbitrary $d \in [n] \setminus \{\dim(e) \mid e \in M\}$. If $|M^d_L| = |M^d_R| = 1$, we are done. Otherwise it must be the case that $M = M^d_i$ for some $i \in \{L, R\}$ and hence by the induction hypothesis, there exists $d' \in [n] \setminus \{d\}$ such that $|M^{d'}_L| = |(M^d_i)^{d'}_L| = 1 = |(M^d_i)^{d'}_R| = |M^{d'}_R|$ as required. □

4 Half-layers

A matching M in Q_n, $n \geq 2$, is called a *half-layer in Q_n* if M consists exactly of all edges of dimension d and parity p for some $d \in [n]$ and $p \in \{0, 1\}$. Observe that each layer is partitioned into two half-layers, each of size 2^{n-2}. Moreover, if a perfect matching of Q_n contains a half-layer, then it actually is a layer of Q_n.

We shall see in the next section that half-layers are obstacles that cannot be overcome. More precisely, we shall see in Lemma 5.1 that there is no Hamiltonian cycle faulting a half-layer. The need to treat half-layers as a special case motivates the following definitions.

Given a matching M in Q_n, we call $d \in [n]$

- a *main dimension of M in Q_n* if M^d contains a half-layer in Q_n;
- a *splitting dimension for M* if M^d_i does not contain a half-layer in Q^d_i for both $i \in \{L, R\}$.

Now we show that the main dimension is uniquely determined.

Proposition 4.1. *Let d be a main dimension of a matching M in Q_n, $n \geq 2$. Then,*

(i) M contains only edges of dimension d. Consequently, d is the only main dimension of M.

(ii) d is the unique main dimension of M^a_i in Q^a_i for every $a \in [n] \setminus \{d\}$ and $i \in \{L, R\}$ provided $n \geq 3$.

Proof. Since for every edge uv of Q_n we have $p(u) \neq p(v)$, it follows that at least one of u, v must be incident with an edge of M^d. Hence $uv \notin M$ and consequently, $M_L^d = M_R^d = \emptyset$. This, together with the fact that M^d contains a half-layer and therefore must be nonempty, verifies part (i).

To see why (ii) holds, note that if $n \geq 3$, then a half-layer in Q_n contains a half-layer in both Q_L^a and Q_R^a by definition. It follows that d is a main dimension of M_i^a in Q_i^a while the uniqueness follows from part (i). □

A matching M in Q_n is called *aligned* if there exists a dimension $d \in [n]$ such that $|M_L^d| \leq 1$ and $|M_R^d| \leq 1$. Otherwise, it is called *unaligned*. Note that a matching containing a half-layer is always aligned by Proposition 4.1(i). For aligned matchings our problem is easy, as we shall see in the next section. But this is not the case with unaligned matchings. For them, we use the following lemma.

Lemma 4.2. *If M is an unaligned matching in Q_n, $n \geq 4$, then there are at least two splitting dimensions for M.*

Proof. First note that if every $a \in [n]$ is a splitting dimension for M, we are done. Hence we can assume that this is not the case, i. e., there exist distinct $a, b \in [n]$ such that b is the main dimension of M_L^a in Q_L^a.

Case 1: b is not a splitting dimension for M. Then there exist $i \in \{L, R\}$ and $d \in [n] \setminus \{b\}$ such that d is the main dimension of M_i^b. Assume that $d \neq a$. Then Proposition 4.1 (ii) implies that d is also the main dimension of $(M_i^b)_L^a$ in $(Q_i^b)_L^a$. It follows that $(M_L^a)_i^b = (M_i^b)_L^a$ contains a half-layer. On the other hand, as b is the main dimension of M_L^a, Proposition 4.1 (i) reveals that M_L^a contains only edges of dimension b and therefore $(M_L^a)_i^b = \emptyset$. This, however, leads to a contradiction since a half-layer is always nonempty. Hence we conclude that $d = a$.

Now choose an arbitrary $c \in [n] \setminus \{a, b\}$ and $j \in \{L, R\}$ and suppose that M_j^c contains a half-layer in Q_j^c. Using Proposition 4.1 (ii) again, a and b are the main dimensions of $(M_i^b)_j^c$ and $(M_L^a)_j^c$ in $(Q_i^b)_j^c$ and $(Q_L^a)_j^c$, respectively. Since $(M_i^b)_j^c = (M_j^c)_i^b$ and $(M_L^a)_j^c = (M_j^c)_L^a$, part (ii) of Proposition 4.1 implies that a and b are also main dimensions of M_j^c in Q_j^c, while part (i) then reveals that $a = b$, which is a contradiction with our assumption. Hence c must be a splitting dimension for M. It follows that there exist $n - 2 \geq 2$ splitting dimensions in this case.

Case 2: b is a splitting dimension for M. In this case we need to show that there exists another splitting dimension for M, different from b. Assume, by way of contradiction, that this is not the case. Then for any $c \in [n] \setminus \{b\}$ there exists $i \in \{L, R\}$ such that M_i^c contains a half-layer. Recall that b is the main dimension of M_L^a in Q_L^a. If $c \neq a$, then Proposition 4.1 (ii) implies that b is the main dimension of $(M_L^a)_i^c = (M_i^c)_L^a$ in $(Q_L^a)_i^c$ and therefore b is also the main dimension of M_i^c in Q_i^c. Hence we conclude that for any $c \in [n] \setminus \{b\}$ there exists $i \in \{L, R\}$ such that b is the main dimension of M_i^c in Q_i^c.

Since M is an unaligned matching, there must exist $i \in \{L, R\}$ such that M_i^b contains two distinct edges e_1, e_2. Note that both $\dim(e_1)$ and $\dim(e_2)$ are different from b. By Proposition 3.5 there must be a dimension $d \in [n] \setminus \{b\}$ such that $|M_j^d \cap \{e_1, e_2\}| = 1$ for each $j \in \{L, R\}$. But then each M_j^d contains an edge of dimension distinct from b, which by Proposition 4.1 (i) means that b cannot be the main dimension of M_j^d, contrary to the conclusion of the previous paragraph. □

5 Admissible configurations

A nonempty matching M is called an *almost-layer* (*of dimension* d) in Q_n if there exist vertices $u, v \in Q_L^d$ of different parities such that M consists of all edges of dimension d except uu^d and vv^d. Pairs $\{u, v\}$ and $\{u^d, v^d\}$ are called the *free pairs* of M.

The next lemma gives necessary conditions for a Gray code faulting a given matching.

Lemma 5.1. *Let M be a matching and u, v vertices of Q_n. If there exists a Hamiltonian path in $Q_n - M$ between u and v, then*

- (i) $p(u) \neq p(v)$;
- (ii) *for every $d \in \Delta(u, v)$ there exists an edge $e \in E(Q_n) \setminus M$ of dimension d such that $d(u, e)$ is odd;*
- (iii) *for every $d \in [n] \setminus \Delta(u, v)$ there exist two edges in $E(Q_n) \setminus M$ of dimension d of different parities;*
- (iv) *M contains no almost-layer with a free pair $\{u, v\}$.*

Proof. Claims (i)–(iii) follow from the facts that in a hypercube the number of vertices of odd parity equals the number of vertices of even parity, and that the parity of vertices on every path alternates. To verify (iv), observe that if M is an almost-layer with a free pair $\{u, v\} \subseteq V(Q_L^d)$, then every path P_{uv} in $Q_n - M$ avoids either all vertices of $V(Q_L^d) \setminus \{u, v\}$, or all vertices of $V(Q_R^d)$, and therefore cannot be Hamiltonian. □

We say that a triple (u, v, M) is a *configuration* in Q_n if M is a matching in Q_n and u and v are vertices of different parities, i. e., condition (i) of

Lemma 5.1 holds. Since we do not need to distinguish between vertices u and v, we consider the configurations (u, v, M) and (v, u, M) to coincide. If, moreover, conditions *(ii)*–*(iv)* of Lemma 5.1 hold, we say that the configuration (u, v, M) is *admissible*.

The conditions *(ii)* and *(iii)* of Lemma 5.1 can be unified into one condition using a slight modification of the concept of a half-layer. Note that since a half-layer contains only edges of the same parity, the distances of a given vertex to edges of a given half-layer must also have the same parity. If the distance of a vertex u to edges of a half-layer M in Q_n is odd, we say that M is an *odd half-layer* for the vertex u.

Proposition 5.2. *A configuration (u, v, M) is admissible if and only if the following holds*

(i) M contains no odd half-layer for u or v; and

(ii) M contains no almost-layer with a free pair $\{u, v\}$.

The next lemma shows that these conditions are also sufficient for the existence of a Hamiltonian path in Q_n, $n \geq 4$, provided that M is aligned.

Lemma 5.3. *Let M be an aligned matching, and let u and v be vertices of Q_n, $n \geq 4$. If (u, v, M) is an admissible configuration, then $Q_n - M$ contains a Hamiltonian path between u and v.*

Proof. Since M is an aligned matching, there exists $d \in [n]$ such that $|M_L^d| \leq 1$ and $|M_R^d| \leq 1$. Without a loss of generality assume that $u \in V(Q_L^d)$ and consider the following two cases:

Case 1: $v \in V(Q_R^d)$. Since (u, v, M) is an admissible configuration, there exists a vertex $w \in V(Q_L^d)$ such that $p(u) \neq p(w)$ and $ww^d \notin M$. Note that then also $p(w^d) \neq p(v)$. Now, by Lemma 3.2, there exist Hamiltonian paths P_{uw} of $Q_L^d - M_L^d$ and $P_{w^d v}$ of $Q_R^d - M_R^d$. Then $P_{uv} := P_{uw} + P_{w^d v}$ is the required Hamiltonian path.

Case 2: $v \in V(Q_L^d)$. Here, we consider the following subcases.

Subcase 2.1: $M_L^d = \emptyset$. Since (u, v, M) is an admissible configuration, there are vertices $x, y \in V(Q_L^d)$ such that $p(x) \neq p(y)$, $xx^d \notin M$, $yy^d \notin M$, and $\{u, v\} \neq \{x, y\}$. Apply Lemma 3.2 to obtain a Hamiltonian path $P_{x^d y^d}$ of $Q_R^d - M_R^d$. Note that as $p(u) \neq p(v)$ and $p(x) \neq p(y)$, we can assume that also $p(u) \neq p(x)$ and $p(v) \neq p(y)$, interchanging x and y if necessary.

We conclude the subcase with the following construction. If $\{u, v\} \cap \{x, y\} = \emptyset$, then apply Lemma 3.3 to obtain spanning paths P_{ux} and P_{vy} of Q_L^d and put $P_{uv} := P_{ux} + P_{x^d y^d} + P_{yv}$. Otherwise it must be the case that $|\{u, v\} \cap \{x, y\}| = 1$. Assume without a loss of generality that $v = x$ and

$u \neq y$. Then apply Lemma 3.1 to obtain a Hamiltonian path P_{uy} of $Q_L^d - v$ and put $P_{uv} := P_{uy} + P_{y^d x^d} + P_{xx}$. In both cases we obtained the required Hamiltonian path.

Subcase 2.2: $M_L^d = \{e\}$. If $e \neq uv$, then put $xy = e$ and apply the construction of Subcase 2.1. So, we may assume that $e = uv$. Since (u, v, M) is an admissible configuration, there is a vertex $y \in V(Q_L^d) \setminus \{u, v\}$ such that $yy^d \notin M$. Assume without a loss of generality that $p(u) \neq p(y)$. Finally, put $x = v$ and apply the construction of Subcase 2.1 in order to conclude this case. □

Note that the above lemma does not hold for $n = 3$, see e. g. configurations (e), (f) or (g) on Figure 1. The case of small dimensions is analyzed separately in the next two sections.

6 Gray codes in Q_3

A configuration (u, v, M) in Q_n is *good* if there is a Hamiltonian path of $Q_n - M$ between u and v. Otherwise it is *bad*. In this terminology, our aim is to characterize all good configurations.

A configuration (u, v, M_1) *contains* a configuration (u, v, M_2) if $M_1 \supseteq M_2$. A bad configuration is *minimal* if it does not contain any other bad configuration. An admissible configuration is *maximal* if it is not contained in any other admissible configuration. Similarly, a *matching* is *maximal* if it is not contained in any other matching.

We start with the list of all minimal bad configurations in Q_3. Note that an electronic version of this paper contains colored figures.

Proposition 6.1. *Let u and v be vertices of different parity and let M be a matching in Q_3. Then there is a Hamiltonian path of $Q_3 - M$ between u and v unless (u, v, M) contains one of the bad configurations depicted on Figure 1.*

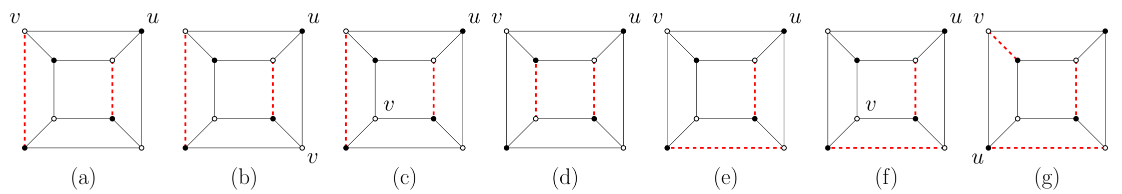

Figure 1: The minimal bad configurations in Q_3.

Proof. First observe that each configuration (u, v, M) on Figure 1 is bad. Note that M is presented with dashed lines. Moreover, observe that there are no two configuration on Figure 1 such that the first one contains the other one.

Now we show that the configuration (u, v, M) is either good or contains one of the bad configurations on Figure 1. The hypercube Q_3 has 10 non-isomorphic matchings, including the empty one, depicted on Figure 2(1)-(10). Note that in this figure the edges of the matchings are presented with dashed lines. So M is one of the following cases on Figure 2.

(1), (2): In these cases M has at most one edge, thus the configuration (u, v, M) is good by Lemma 3.2.

(3), (6): Observe that if $u \in \{u_1, u_2\}$ and $v \in \{v_1, v_2\}$, then the configuration (u, v, M) is good. Otherwise it contains a configuration of type (a), (b), or (c) on Figure 1.

(4): Observe that the configuration (u,v,M) is good unless it equals (u_1,v_1,M) or (u_2, v_2, M) which are of type (d) on Figure 1.

(5): Observe that the configuration (u,v,M) is good unless it equals (u_1,v_1,M), (u_1, v_2, M), or (u_2, v_2, M) which are of type (e) or (f) on Figure 1.

(7): In this case $Q_3 - M - xy$ is a Hamiltonian cycle. Thus the configuration (u, v, M) is good if uv is an edge on this cycle. Furthermore, if the configuration (u, v, M) equals (u_1, v_1, M) or (u_2, v_2, M), observe that it is also good. Otherwise it contains a configuration of type (e) or (f) on Figure 1.

(8): Observe that if $uv \in M$, then the configuration (u, v, M) is good. Furthermore, it is also good if $d(u, v) = 3$ and $\{u, v\} \neq \{u_1, v_1\}$. Otherwise it contains a configuration of type (e), (f), or (g) on Figure 1.

(9): In this case $Q_3 - M$ is disconnected so (u, v, M) is a bad configuration. Observe that it contains a configuration of type (a), (b), or (c) on Figure 1.

(10): In the last case $Q_3 - M$ is a Hamiltonian cycle. Thus the configuration (u, v, M) is good if uv is an edge on this cycle. Otherwise it contains a configuration of type (e) or (f) on Figure 1.

This completes the analysis of all cases and establishes the lemma. □

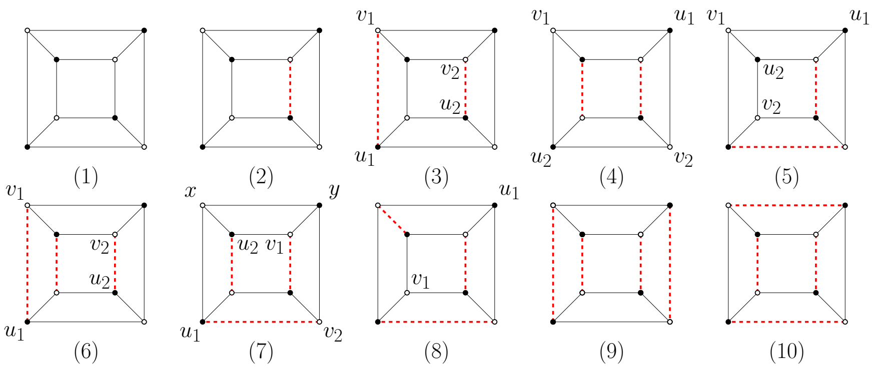

Figure 2: All matchings in Q_3 up to isomorphism.

The following two lemmas will be useful for inquiry of Q_4 in the next section. The first lemma applies typically for configurations (u, v, M) such that both u and v belong to a same subcube Q_L^d or Q_R^d which contains at most 2 edges of M.

Lemma 6.2. *Let u and v be vertices of different parity and let e_1 and e_2 be distinct edges in Q_3. Then there is a Hamiltonian path of Q_3 between u and v that*

(i) contains exactly one of e_1 and e_2; or

(ii) faults both e_1, e_2 but contains an edge xy such that $x \in e_1$ and $y \in e_2$.

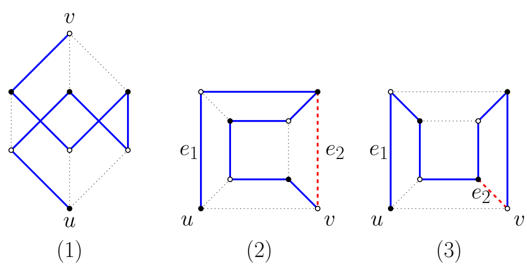

Figure 3: An illustration for Lemma 6.2.

Proof. Suppose first that u and v are at distance 3. From Figure 3(1) it is obvious that all Hamiltonian paths between u and v are isomorphic. Moreover, each of them contains one edge of the upper level, one edge of the lower level, and five edges of the middle level. Considering few cases, regarding the levels that contain e_1 and e_2, it is easy to find a suitable path.

Suppose now that uv is an edge of Q_3. We distinguish three cases:

Case 1: Edges e_1 and e_2 are adjacent at some vertex x. If $x = u$, then assume that $e_1 = uy \neq uv$ and apply Lemma 3.1 to find a Hamiltonian path of $Q_3 - u$ between y and v. Extending this path with the edge $e_1 = uy$ we obtain a Hamiltonian path of Q_3 between u and v that contains e_1 and faults e_2. If $x = v$, then proceed similarly. Otherwise we have $x \notin \{u, v\}$. The Hamiltonian path between u and v that faults e_1 from Lemma 3.2 contains e_2 since it passes through the vertex x.

Case 2: *Edges e_1 and e_2 are at distance 1.* If the edge uv is adjacent to both e_1 and e_2, then we have two non-isomorphic cases on Figures 3(2)-(3) with desired Hamiltonian paths. Otherwise, we may assume that $x \in e_1$, $y \in e_2$, and $y \notin \{u, v\}$ for some edge xy. The Hamiltonian path between u and v that faults e_1 from Lemma 3.2 contains e_2 or xy since it passes through the vertex y.

Case 3: *Edges e_1 and e_2 are at distance 2.* If $e_1 = uv$, apply Corollary 3.4 to find a Hamiltonian path between u and v that contains edge e_2. If $e_2 = uv$, proceed similarly. Otherwise the configuration $(u, v, \{e_1, e_2\})$ is bad since it is of type (a), (b), or (c) on Figure 1. Thus the Hamiltonian path between u and v that faults e_1 from Lemma 3.2 contains e_2. This completes the analysis of all cases. $\square$

The second lemma applies for example when $M^d_R = \{ab, cd\}$ as on Figure 4 and Q^d_L contains a path between vertices b^d and c^d. The extra edge bc represents this path.

Lemma 6.3. *Let a, b, c, d, e, f be vertices of Q_3 as depicted on Figure 4 and $u, v \in V(Q_3)$ be any two vertices of different parity. Then $Q_3 + bc$ has a Hamiltonian path between u and v that contains the edge bc and faults $\{ab, cd\}$ unless $\{u, v\} = \{b, c\}$ or $\{u, v\} = \{e, f\}$.*

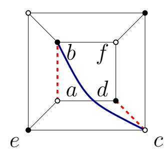

Figure 4: The configuration of Lemma 6.3.

Proof. By exhausting all 14 possibilities for $\{u, v\}$, as it is presented on Figure 5, one can verify the validity of the lemma. $\square$

7 Gray codes in Q_4

If M is aligned in Q_4, we can use Lemma 5.3 which holds for every $n \geq 4$. Thus, it remains to consider admissible configurations (u, v, M) only with

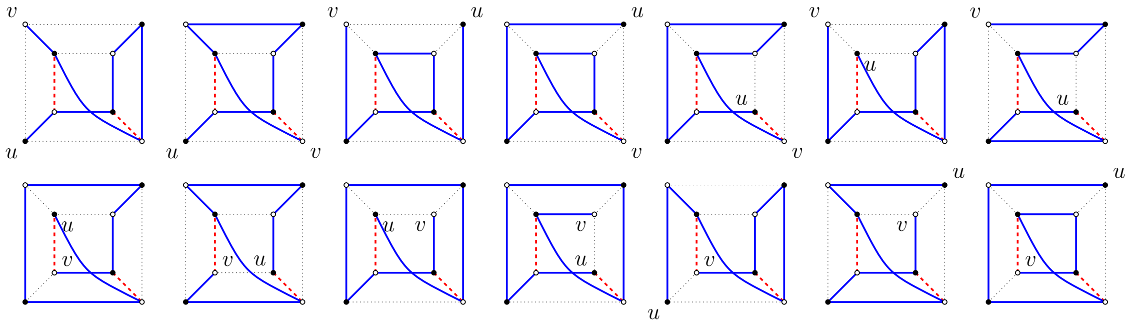

Figure 5: Hamiltonian paths in Lemma 6.3 for each pair $\{u, v\}$ of endvertices.

an unaligned matching M. We start with the list of all maximal unaligned matchings in Q_4.

Let $m(M)$ denote the maximum number of edges of M of the same dimension, i. e., $m(M) = \max_{d\in[n]} |M^d|$.

Lemma 7.1. *Let M be a maximal unaligned matching in Q_4. Then $m(M) \in \{2, 3, 4\}$ and M is isomorphic to one of the matchings on Figures 6-8.*

Proof. Let $d \in [4]$ be a dimension with the maximal number of edges of M, i. e., $|M^d| = m(M)$. Since M is unaligned, we have $|M^d_L| \geq 2$ or $|M^d_R| \geq 2$, say $|M^d_L| \geq 2$. Thus among eight vertices of Q^d_L there are at most four vertices that are not matched by M^d_L. Hence $m(M) \leq 4$. On the other hand, every maximal matching in Q_4 contains at least 6 edges [6], hence at least 2 edges from M are of a same dimension. This implies that $m(M) \geq 2$. So $m(M) \in \{2, 3, 4\}$.

Let G_L and G_R be the graphs obtained from Q^d_L and Q^d_R, respectively, by removing all vertices incident with M^d. Note that G_L and G_R are isomorphic, so let us write simply G when it is clear from the context for which of the graphs we speak. Since M is maximal, both M^d_L and M^d_R are maximal matchings of G.

To show that M is one of the matchings on Figures 6-8, we first consider the set M^d and then the combination of sets M^d_L and M^d_R. The matchings M^d, M^d_L and M^d_R on the following figures are represented with dashed lines, the edges of the graph G are bold, and the remaining edges are dotted.

Note that the graphs G_R and G are flipped on the following figures, so the matchings M^d_R and the maximal matchings of G are also flipped.

Case 1: $m(M) = 4$. Since M is unaligned, there are three non-isomorphic combinations of four edges in M^d; see Figure 9. We consider now three possibilities.

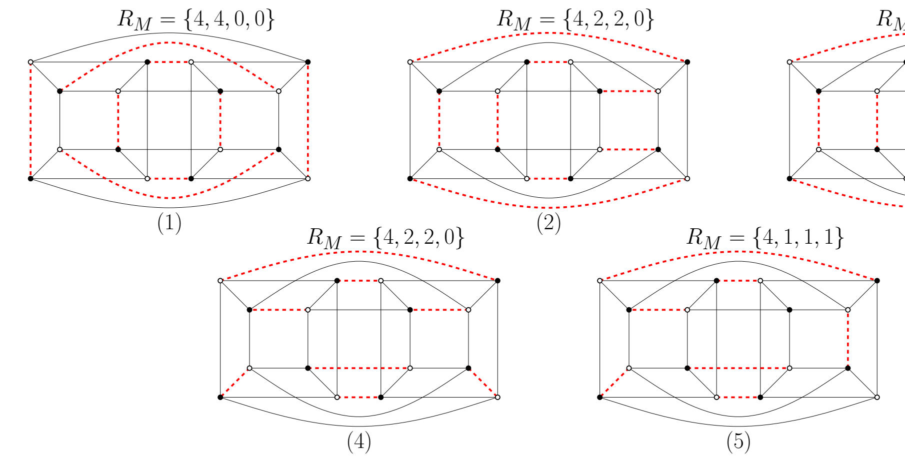

Figure 6: All maximal unaligned matchings M of Q_4 with $m(M) = 4$.

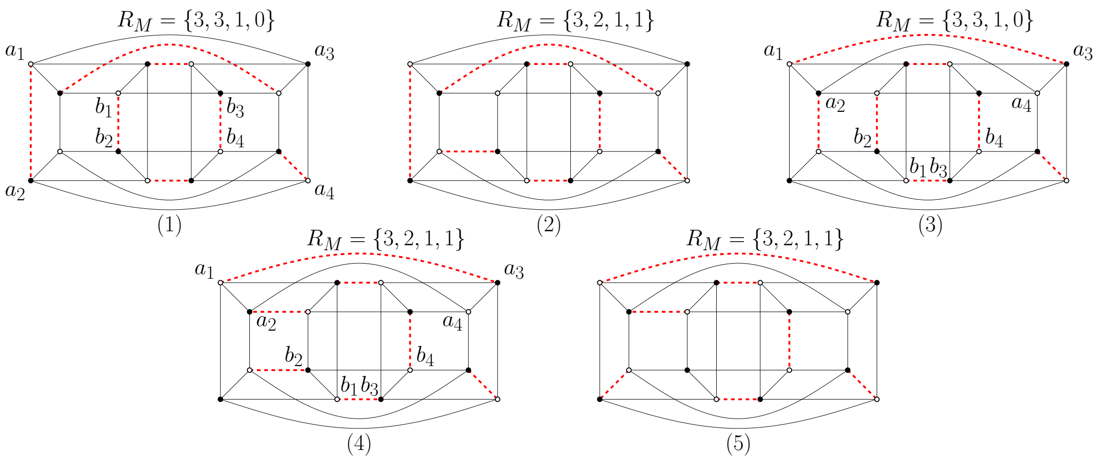

Figure 7: All maximal unaligned matchings M of Q_4 with $m(M) = 3$.

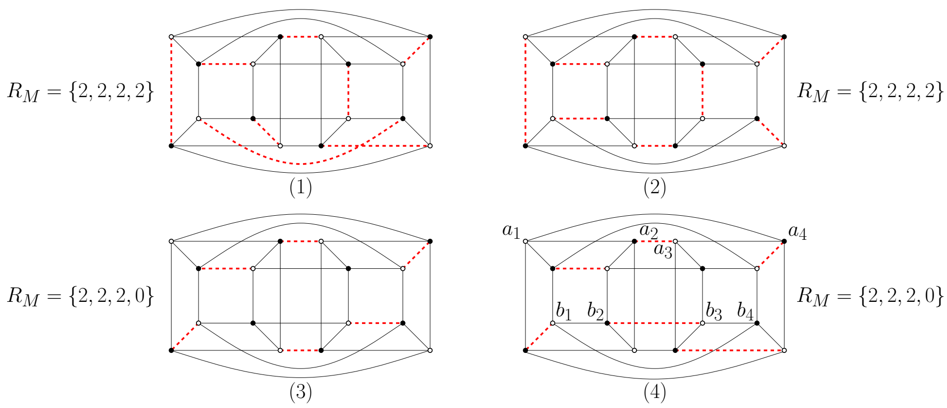

Figure 8: All maximal unaligned matchings M of Q_4 with $m(M) = 2$.

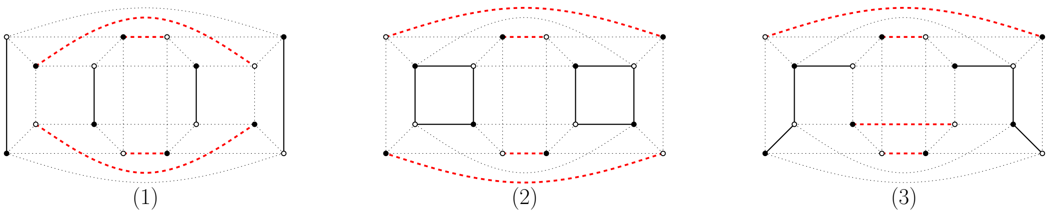

Figure 9: All sets M^d with 4 edges in a maximal unaligned matching M of Q_4.

Subcase 1.1: M^d is as on Figure 9(1). Obviously, there is only one maximal matching of G. Thus M^d_L and M^d_R are uniquely determined and the matching M is as on Figure 6(1).

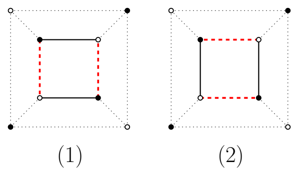

Figure 10: All maximal matchings of G in Subcase 1.2.

Subcase 1.2: M^d is as on Figure 9(2). There are two maximal matchings of G depicted on Figure 10. By symmetry, we may assume that M^d_L is as on Figure 10(1). Now, if M^d_R is also as on Figure 10(1), the matching M corresponds to Figure 6(3). And, if M^d_R is as on Figure 10(2), then M corresponds to Figure 6(2).

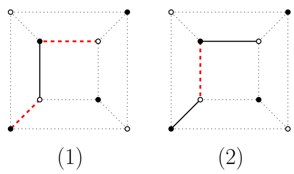

Figure 11: All maximal matchings of G in Subcase 1.3.

Subcase 1.3: M^d is as on Figure 9(3). There are two maximal matchings of G depicted on Figure 11. By the maximality of M, it follows that M^d_L and M^d_R are not both as on Figure 11(2). Now, if both of them are as on Figure 11(1), we infer Figure 6(4). Finally, if one of them is as on Figure 11(1), and the other one as on Figure 11(2), we infer Figure 6(5).

Case 2: $m(M) = 3$. Since M is unaligned, there are two non-isomorphic combinations of three edges in M^d which are depicted on Figure 12. Consider now these two cases separately.

Subcase 2.1: M^d is as on Figure 12(1). There are three maximal matchings of G depicted on Figure 13. By symmetry, we may assume that if M^d_L

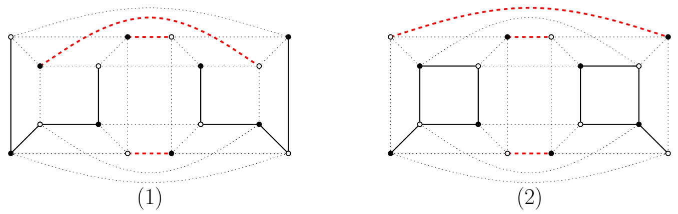

Figure 12: All sets M^d with 3 edges in a maximal unaligned matching M of Q_4.

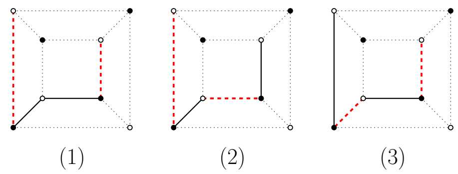

Figure 13: All maximal matchings of G in Subcase 2.1.

is as on Figure 13(x), then M_R^d is as on Figure 13(y) where $1 \leq x \leq y \leq 3$. Moreover, since M is a maximal matching, we infer $x \neq y$. We have the following possibilities for the pair (x, y):

$(1, 3)$: The matching M is on Figure 7(1);

$(2, 3)$: The matching M is on Figure 7(2);

$(1, 2)$: This case is isomorphic to the case $(1, 3)$. Indeed, consider the isomorphism of Q_4 that interchanges vertices a_i and b_i on Figure 7(1) for every $i \in [4]$ and fixes all other vertices.

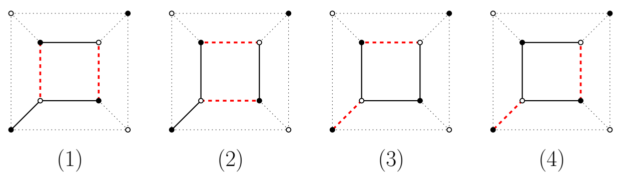

Figure 14: All maximal matchings of G in Subcase 2.2.

Subcase 2.2: M^d is as on Figure 12(2). There are four maximal matchings of G depicted on Figure 14. By symmetry, we may assume that if M_L^d is as on Figure 14(x), then M_R^d is as on Figure 14(y) where $1 \leq x \leq y \leq 4$. Moreover, since M is a maximal matching, it follows that $x \neq y$. Thus, we have the following possibilities for the pair (x, y):

$(1, 4)$: The matching M is on Figure 7(3);

$(2,4)$: The matching M is on Figure 7(4);

$(3,4)$: The matching M is on Figure 7(5);

$(1,2)$: This case does not occur since M is a maximal matching;

$(1,3)$ and $(2,3)$: These cases are isomorphic to the cases $(2,4)$ and $(1,4)$, respectively. Indeed, consider the isomorphism of Q_4 that interchanges vertices a_i and b_i on Figures 7(4) and 7(3) for every $i \in [4]$ and that fixes all other vertices.

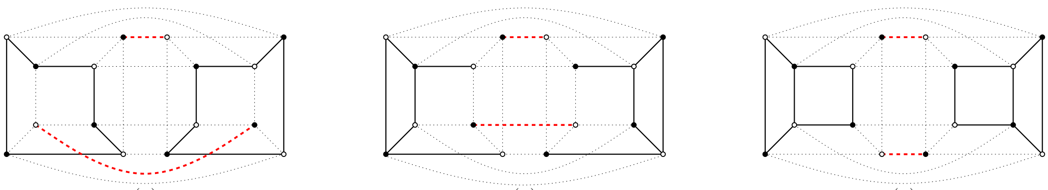

Figure 15: All sets M^d with 2 edges in a maximal unaligned matching M of Q_4.

Case 3: $m(M) = 2$. Let $k(M)$ be the minimal distance of distinct edges of M of a same dimension, i. e.,

$$k(M) = \min\{d(e_1, e_2) \mid e_1, e_2 \in M^i \text{ for some } i \in [4] \text{ and } e_1 \neq e_2\}.$$

We assume that M^d contains edges at distance $k(M)$, otherwise choose d accordingly. For each $k(M) \in [3]$ we have M^d on Figure 15.

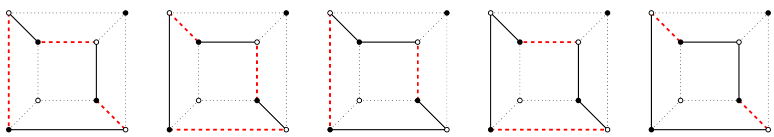

Figure 16: All maximal matchings of G in Subcase 3.1.

Subcase 3.1: $k(M) = 3$. Thus M^d is as on Figure 15(1). There are five maximal matchings of G depicted on Figure 16. Observe that neither M_L^d nor M_R^d is as on Figures 16(3)-(5) since $k(M) = 3$. Moreover, M_d^L and M_d^R are not both as on Figure 16(1), or both as on Figure 16(2) since $k(M) = 3$. Thus, one of them is as on Figure 16(1), and the other one as on Figure 16(2). Hence M corresponds to Figure 8(1).

Subcase 3.2: $k(M) = 2$. Thus M^d is as on Figure 15(2). There are seven maximal matchings of G depicted on Figure 17. Observe that neither M_L^d

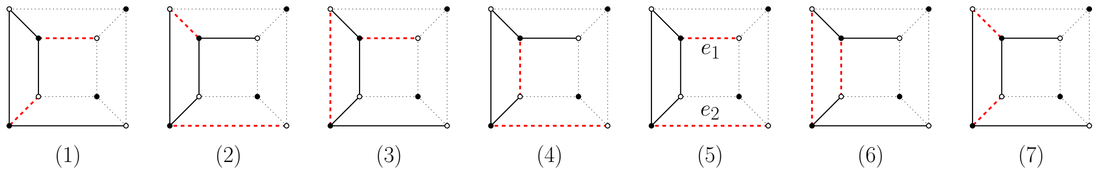

Figure 17: All maximal matchings of G in Subcase 3.2.

nor M_R^d is as on Figures 17(6)-(7) since $k(M) = 2$. Furthermore, neither M_L^d nor M_R^d is as on Figure 17(5) since all matchings of G on Figures 17(1)-(5) contain at least one edge of e_1 or e_2 and $k(M) = 2$. By symmetry, we may assume that if M_L^d is as on Figure 17(x), then M_R^d is as on Figure 17(y) where $1 \le x \le y \le 4$. Moreover, $x \ne y$ since M is a maximal matching. We have the following possibilities for the pair (x, y):

$(1, 2)$: The matching M is on Figure 8(4);

$(3, 4)$: This case is isomorphic to the case $(1, 2)$. Indeed, consider the isomorphism of Q_4 that interchanges vertices a_i and b_i on Figure 8(4) for every $i \in [4]$ and fixes all other vertices.

$(1, 3)$, $(1, 4)$, $(2, 3)$, and $(2, 4)$: These cases do not occur since M is a maximal matching.

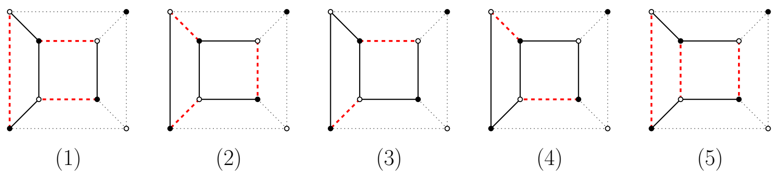

Figure 18: All maximal matchings of G in Subcase 3.3.

Subcase 3.3: $k(M) = 1$. Thus M^d is as on Figure 15(3). There are five maximal matchings of G depicted on Figure 18. Observe that neither M_L^d nor M_R^d is as on Figure 18(5) since $m(M) = 2$. By symmetry, we may assume that if M_L^d is as on Figure 18(x), then M_R^d is as on Figure 18(y) where $1 \le x \le y \le 4$. We have the following possibilities for the pair (x, y):

$(1, 2)$: The matching M is on Figure 8(2);

$(3, 4)$: The matching M is on Figure 8(3);

$(1, 1)$, $(1, 3)$, $(1, 4)$, $(2, 2)$, $(2, 3)$, and $(2, 4)$: These cases do not occur since $m(M) = 2$;

$(3,3)$ and $(4,4)$: These cases do not occur since M is a maximal matching.

Note that the choice of the dimension d was not deterministic. Therefore, to conclude the proof, we need to verify that the matchings on Figures 6-8 are non-isomorphic.

Define the *mark* of M to be the collection $R_M = \{|M^d|\}_{d\in[n]}$. Clearly, matchings with different marks are not isomorphic. So we need to examine only matchings with different marks on Figures 6-8. This can be easily verified by the reader. □

Lemma 7.2. *Let (u,v,M) be a maximal admissible configuration in Q_n with an unaligned matching M. Then M is maximal with respect to inclusion.*

Proof. Suppose on the contrary that M is not maximal, i. e., $M \cup \{e\}$ is a matching for some edge $e \in E(Q_n) \setminus M$. Let d be the dimension of e. Since M is unaligned, we have $|M_L^d| \geq 2$ or $|M_R^d| \geq 2$, so we may assume that M_L^d contains edges x_1x_2 and x_3x_4. Let $e_i = x_ix_i^d$ for $i \in [4]$.

Since $d(e_1,e_2)$ is odd, $d(e_3,e_4)$ is odd, e_1,e_2,e_3,e_4 are not in $M \cup \{e\}$, and (u,v,M) is admissible, it follows from the definition that $(u,v,M\cup\{e\})$ is also admissible. But this contradicts the assumption that (u,v,M) is a maximal admissible configuration. □

Now we are ready to find a Hamiltonian path of $Q_4 - M$ between u and v if (u,v,M) is admissible configuration. The following lemma serves as a basis of induction in the next section.

Lemma 7.3. *Let M be a matching in Q_4 and $u,v \in V(Q_4)$. If (u,v,M) is admissible, then there is a Hamiltonian path of $Q_4 - M$ between u and v.*

Proof. If M is aligned, we can use Lemma 5.3 which holds for every $n \geq 4$. Thus, it remains to consider admissible configurations (u,v,M) only with unaligned matching M.

Assume that (u,v,M) is a maximal admissible configuration, otherwise introduce new edges in M. Notice that with this procedure M stays unaligned. By Lemma 7.2, the matching M is maximal. Thus $m(M) \in \{2,3,4\}$ and M corresponds to one of the matchings on Figures 6-8 by Lemma 7.1.

Let $d \in [4]$ be a dimension with the maximal number of edges from M, i. e., $|M^d| = m(M)$. Now, as d is fixed, for the sake of simplicity, let us omit d in the notions of Q_L^d, Q_R^d, M_L^d, and M_R^d. Furthermore, for $x \in V(Q_R)$ and $y \in V(Q_L)$ let x_L and y_R denote their corresponding vertices in the opposite subcube, i. e., $x_L = x^d$ and $y_R = y^d$. We say that vertices u and v are *separated* if $u \in V(Q_L)$ and $v \in V(Q_R)$ or vice versa.

In what follows, we distinguish three cases regarding $m(M)$. In each case we consider good/bad configurations in Q_L and Q_R. Then, regarding the position of vertices u and v, we look separately in Q_L and Q_R for paths that can be combined into a Hamiltonian path of $Q_4 - M$ between u and v.

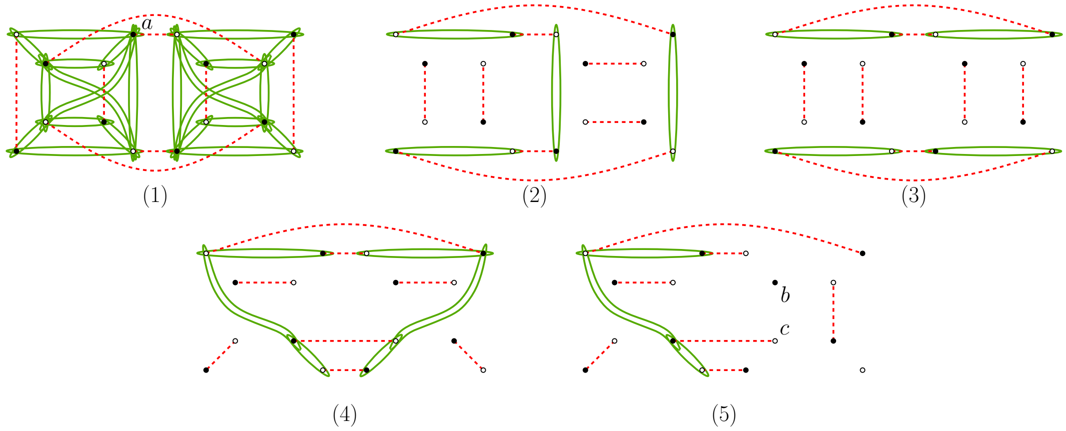

Figure 19: All bad configurations in Q_L or Q_R for each maximal unaligned matching M in Q_4 with $m(M) = 4$.

Case 1: $m(M) = 4$. Applying Proposition 6.1, observe that Figure 19(1)-(5) shows all bad configurations in Q_L or Q_R for each of the five maximal unaligned matchings M with $m(M) = 4$ depicted on Figure 6, respectively. Vertices u and v that form a bad configuration (u, v, M_L) in Q_L or (u, v, M_R) in Q_R are encompassed by a solid elliptic curve, the matching M is represented by dashed edges, whereas other edges of Q_4 are omitted for the sake of clarity. Thus e. g. on Figure 19(2) we have two bad configurations in Q_L and two bad configurations in Q_R.

Subcase 1.1: Vertices u and v are separated. Assume that $u \in V(Q_L)$ and $v \in V(Q_R)$. We distinguish the following two possibilities.

1.1.1: M is as on Figure 19(1) and u or v is incident with M^d. By symmetry, we may assume that $u = a$. Observe on Figure 20 that the statement of the lemma holds for each $v \in V(Q_R)$ of different parity than u. The Hamiltonian paths are presented by bold edges.

1.1.2: M is as on Figure 19(2)-(5) or none of u and v is incident with M^d. Let $w \in V(Q_L)$ be a vertex of different parity than u such that it is incident with M_L. Then, v and w^R are also of different parities. Observe on Figure 19(2)-(5) that the vertices w and w^R are in no bad configuration in Q_L and Q_R, respectively. Thus the configurations (u, w, M_L) and (v, w_R, M_R) are good.

Hence there is a Hamiltonian path P_{uw} of $Q_L - M_L$, and a Hamiltonian path $P_{w_R v}$ of $Q_R - M_R$. The desired Hamiltonian path of $Q_4 - M$ is $P_{uw} + P_{w_R v}$.

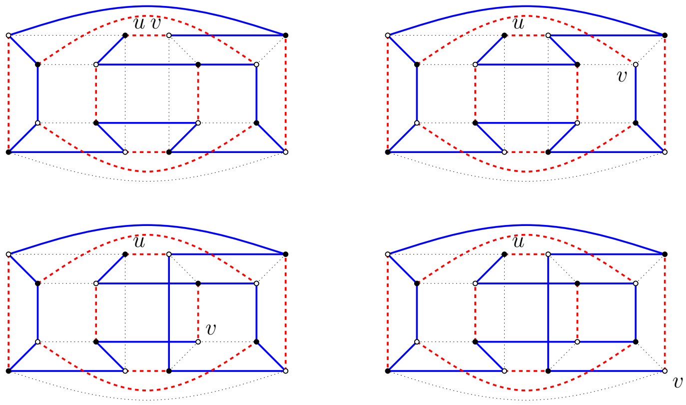

Figure 20: Hamiltonian paths for Subcase 1.1.1.

Subcase 1.2: Both vertices u and v are in Q_L. Note that M_L is of size 2, thus denote by e_1, e_2 these two edges. By Lemma 6.2, there is a Hamiltonian path P_{uv} that either contains precisely one of e_1 and e_2, or faults both e_1, e_2 but contains an edge adjacent to both e_1, e_2. Let $xy \in E(Q_L)$ be that contained edge. Observe on Figure 19(1)-(5) that there is no bad configuration (w, z, M_R) in Q_R such that w_L and z_L are both incident with e_1 or e_2. Thus the configuration (x_R, y_R, M_R) is good.

Hence there is a Hamiltonian path $P_{x_R y_R}$ of $Q_R - M_R$. The desired Hamiltonian path of $Q_4 - M$ is $P_{ux} + P_{x_R y_R} + P_{yv}$ where P_{ux} and P_{yv} are disjoint subpaths of P_{uv}, assuming that x is closer to u on P_{uv} than y.

Subcase 1.3: Both vertices u and v are in Q_R. If M is as on Figure 19(1)-(4), we have symmetry to the previous subcase. However, if M is as on Figure 19(5), then M_R contains only one edge e_1. Then just put $e_2 = bc$ and proceed symmetrically as in Subcase 1.2.

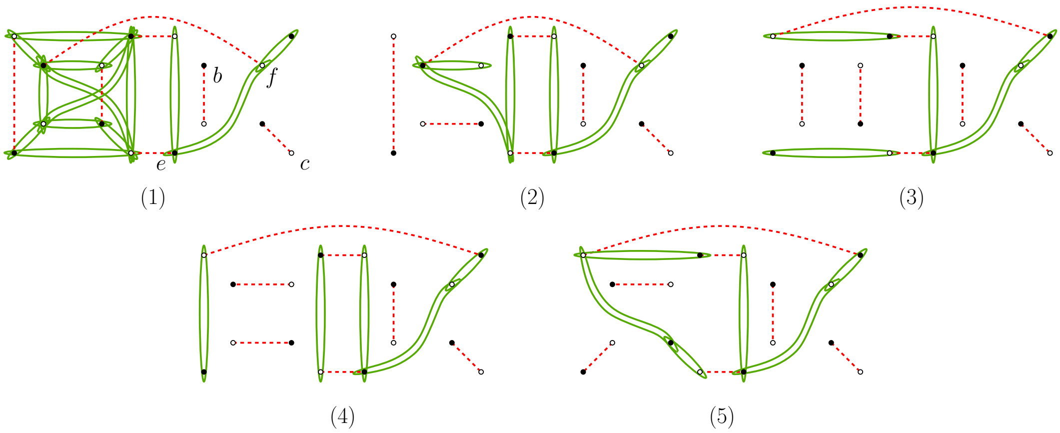

Figure 21: All bad configurations in Q_L or Q_R for each maximal unaligned matching M in Q_4 with $m(M) = 3$.

Case 2: $m(M) = 3$. Applying Proposition 6.1 observe that Figure 21(1)-(5) shows all bad configurations in Q_L or Q_R for each of the five maximal unaligned matchings M with $m(M) = 3$ depicted on Figure 7, respectively. We consider the following possibilities.

Subcase 2.1: Vertices u and v are separated. Again, assume that $u \in V(Q_L)$ and $v \in V(Q_R)$. We distinguish the following two possibilities.

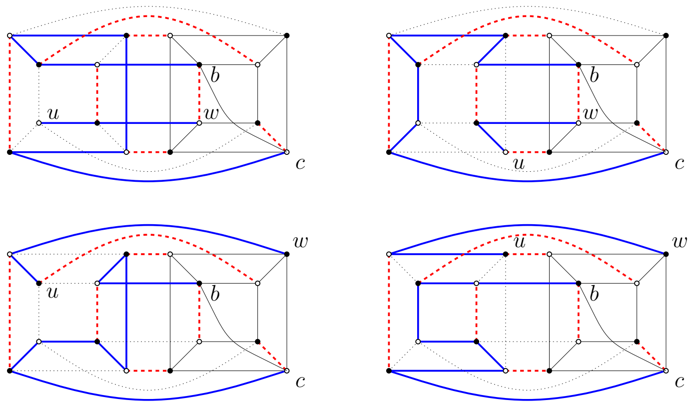

Figure 22: Spanning paths of Q_L for Subcase 2.1.1.

2.1.1: M is as on Figure 21(1) and u is not incident with M_L. Observe on Figure 22 that for each $u \in V(Q_L)$ that is not incident with M_L, there are spanning paths P_{uw_L} and $P_{b_Lc_L}$ of $Q_4 - M_L$. By Lemma 6.3, there is a Hamiltonian path P_{wv} of $Q_R + bc$ that contains the edge bc and faults M_R.

Hence a desired Hamiltonian path of $Q_4 - M$ is $P_{uw_L} + P_{wb} + P_{b_Lc_L} + P_{cv}$ where P_{wb} and P_{cv} are disjoint subpaths of P_{wv}, assuming that b is closer to w on P_{wv} than c.

2.1.2: M is as on Figure 21(2)-(5) or u is incident with M_L. Let $x, y \in V(Q_L)$ be vertices of different parity than u that are incident with M_L. We can always choose such x, y since $|M_L| = 2$. Observe on Figure 21(1) that each bad configuration in Q_L includes at most one vertex incident with M_L. Similarly, observe on Figure 21(2)-(5) that each bad configuration in Q_L includes no vertex incident with M_L. Thus both configurations (u, x, M_L) and (u, y, M_L) are good. Moreover, observe on Figure 21(1)-(5) that for at most one $z \in V(Q_L)$, if z is incident with M_L, then z_R is in a bad configuration in Q_R. Hence at most one of configurations (v, x_R, M_R) and (v, y_R, M_R) is bad. Let $w \in \{x, y\}$ be such that the configuration (v, w_R, M_R) is good.

Therefore, there is a Hamiltonian path P_{uw} of $Q_L - M_L$, and a Hamiltonian path P_{w_Rv} of $Q_R - M_R$. Finally, the desired Hamiltonian path is $P_{uw} + P_{w_Rv}$.

Subcase 2.2: Vertices u and v are both in Q_L or both in Q_R. We distinguish the following two possibilities.

2.2.1: M is as on Figure 21(1) and both u and v are in Q_R. If $\{u, v\} =$

$\{b,c\}$ or $\{u,v\} = \{e,f\}$, observe on Figure 23 that the statement of the lemma holds. Otherwise, by Lemma 6.3, there is a Hamiltonian path P_{uv} of $Q_R + bc$ that contains the edge bc and faults M_R. Observe on Figure 21(1) that the configuration (b_L, c_L, M_L) is good.

Thus there is a Hamiltonian path $P_{b_Lc_L}$ of $Q_L - M_L$. The desired Hamiltonian path of $Q_4 - M$ is $P_{ub} + P_{b_Lc_L} + P_{cv}$ where P_{ub} and P_{cv} are disjoint subpaths of P_{uv}, assuming that b is closer to u on P_{uv} than c.

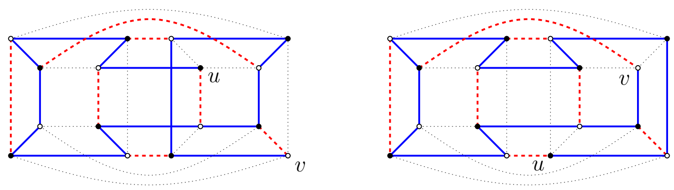

Figure 23: Hamiltonian paths for Subcase 2.2.1.

2.2.2: M is as on Figure 21(1) and both u and v are in Q_L, or M is as on Figure 21(2)-(5). We proceed similarly as in Subcase 1.2. Assume that $u, v \in V(Q_L)$, otherwise interchange the roles of Q_L and Q_R. Let $M_L = \{e_1, e_2\}$. By Lemma 6.2, there is a Hamiltonian path P_{uv} that either contains precisely one of e_1 and e_2, or faults both e_1, e_2 but contains an edge adjacent to both e_1, e_2. Let xy be that contained edge. Observe on Figure 21(1)-(5) that there is no bad configuration (w, z, M_R) in Q_R such that w_L and z_L are both incident with e_1 or e_2. Thus the configuration (x_R, y_R, M_R) is good.

Hence there is a Hamiltonian path $P_{x_Ry_R}$ of $Q_R - M_R$. The desired Hamiltonian path of $Q_4 - M$ is $P_{ux} + P_{x_Ry_R} + P_{yv}$ where P_{ux} and P_{yv} are disjoint subpaths of P_{uv}, assuming that x is closer to u than y on the path P_{uv}.

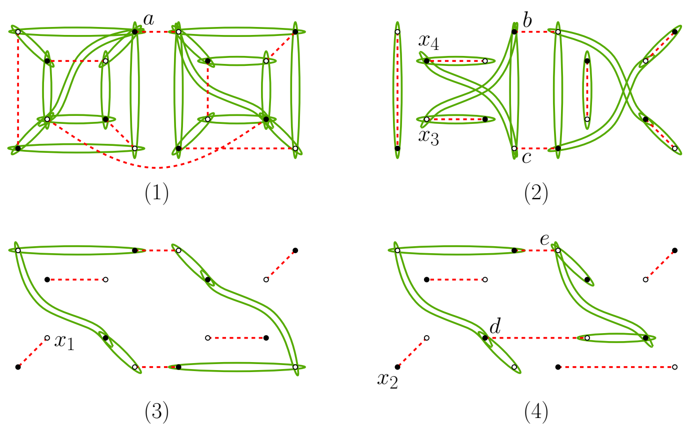

Figure 24: All bad configurations in Q_L or Q_R for each maximal unaligned matching M in Q_4 with $m(M) = 2$.

Case 3: $m(M) = 2$. Applying Proposition 6.1 observe that Figure 24(1)-(4) shows all bad configurations in Q_L or Q_R for each of the four maximal unaligned matchings M with $m(M) = 2$ depicted on Figure 8, respectively.

Subcase 3.1: Vertices u and v are separated. Assume that $u \in V(Q_L)$ and $v \in V(Q_R)$. We distinguish the following three possibilities.

3.1.1: M is as on Figure 24(1) and u or v is incident with M^d. Assume that $u = a$, the other case is symmetric. Observe on Figure 25 that the statement holds for each $v \in V(Q_R)$ of different parity than u.

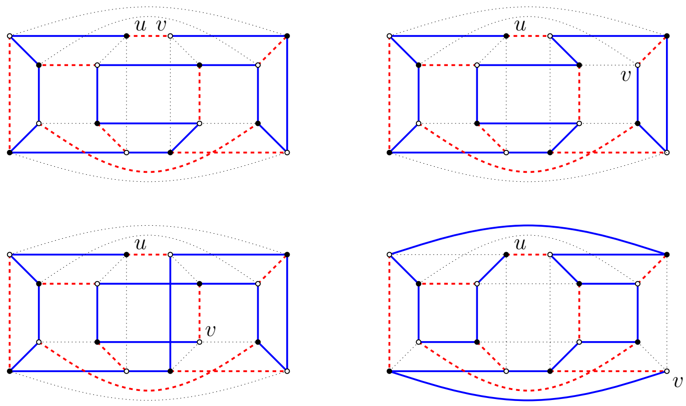

Figure 25: Hamiltonian paths for Subcase 3.1.1.

3.1.2: M is as on Figure 24(4) and $\{u, v\} = \{d, e\}$. Observe on Figure 26 that the statement holds.

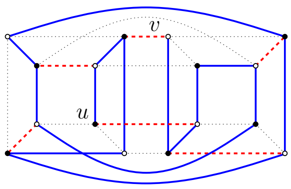

Figure 26: A Hamiltonian path for Subcase 3.1.2.

3.1.3: The remaining possibility; that is, M is as on Figure 24(1) and none of u and v is incident with M^d, or M is as on Figure 24(2)-(4) and $\{u, v\} \neq \{d, e\}$. A vertex $w \in V(Q_L)$ is *free* if it has a different parity than u and is not incident with M^d.

We claim that *there is a free vertex $w \in V(Q_L)$ such that the configurations (u, w, M_L) and (v, w_R, M_R) are good.*

Suppose first that M is as on Figure 24(1)-(3). Observe from the same figure that the configuration (u, x, M_L) is bad for at most one free vertex x. Similarly, the configuration (v, y_R, M_R) is bad for at most one free vertex y. Since there are three free vertices, the claim holds in this case.

Suppose now that M is as on Figure 24(4). If u and d have different parity, then we have two free vertices and both of them satisfy the claim. If u and d have the same parity, we have four free vertices. Observe on Figure 24(4) that the configurations (u, x, M_L) or (v, x_R, M_R) are bad for at most three free vertices x, since $\{u, v\} \neq \{d, e\}$. Hence the claim holds in all cases.

Therefore, there is a Hamiltonian path P_{uw} of $Q_L - M_L$, and a Hamiltonian path $P_{w_R v}$ of $Q_R - M_R$. The desired Hamiltonian path of $Q_4 - M$ is $P_{uw} + P_{w_R v}$.

Subcase 3.2: Vertices u and v are both in Q_L or both in Q_R. Assume that $u, v \in V(Q_L)$, the case $u, v \in V(Q_R)$ is symmetric. We distinguish the following five possibilities.

3.2.1: M is as on Figure 24(1) and the configuration (u, v, M_L) is bad. We have three non-isomorphic cases: both, precisely one, or none of u and v is incident with M_L. Observe on Figure 27 that the statement holds for each of these three possibilities.

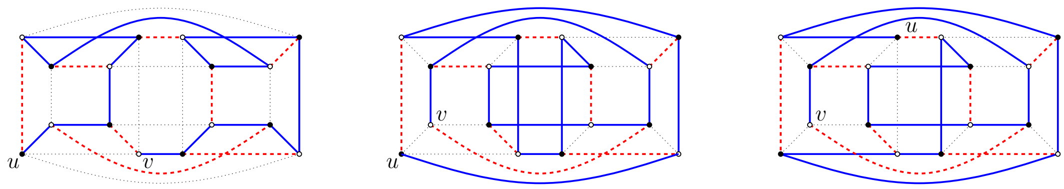

Figure 27: Hamiltonian paths for Subcase 3.2.1.

3.2.2: M is as on Figure 24(1) and the configuration (u, v, M_L) is good. We have two non-isomorphic cases: $uv \in M_L$ and $uv \notin M_L$. Observe on Figure 28 that the statement holds for both of these cases.

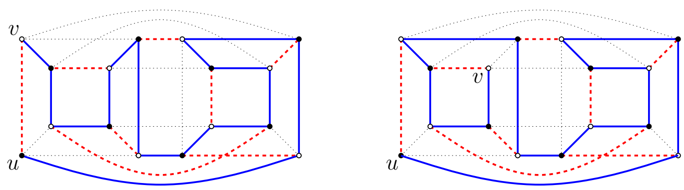

Figure 28: Hamiltonian paths for Subcase 3.2.2.

3.2.3: M is as on Figure 24(2) and $\{u, v\} = \{b, c\}$ or $\{u, v\} = \{x_3, x_4\}$. Observe on Figure 29 that the statement holds.

3.2.4: M is as on Figure 24(2)-(4), the configuration (u, v, M_L) is good, and $\{u, v\} \neq \{x_3, x_4\}$. Then there is a Hamiltonian path P_{uv} of $Q_L - M_L$.

We claim that *P_{uv} contains an edge $xy \in P_{uv}$ such that $xx^R \notin M^d$, $yy^R \notin M^d$, and the configuration (x_R, y_R, M_R) is good.*

First, if M is as on Figure 24(3) or 24(4), put $x = x_1$ or $x = x_2$, respectively, and choose y to be any neighbor of x on P_{uv}. Observe on

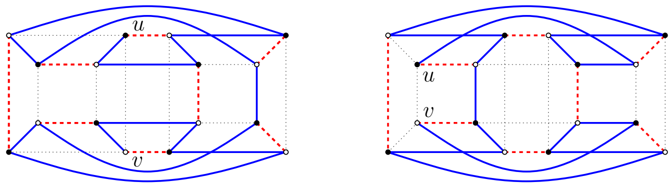

Figure 29: Hamiltonian paths for Subcase 3.2.3.

Figure 24(3)-(4) that x is in no bad configuration in Q_L, also x_R is in no bad configuration in Q_R, $xx^R \notin M^d$, and $yy^R \notin M^d$. Thus the claim holds in this situation.

Now M is as on Figure 24(2). Since $\{u, v\} \neq \{x_3, x_4\}$, we have $x_3 \notin \{u, v\}$ or $x_4 \notin \{u, v\}$, say $x_3 \notin \{u, v\}$. Thus the Hamiltonian path P_{uv} contains the edge x_3x_4. Put $x = x_3$ and $y = x_4$ and observe on Figure 24(2) that the claim is established.

Hence, by the claim, there is a Hamiltonian path $P_{x_Ry_R}$ of $Q_R - M_R$. Finally, the desired Hamiltonian path of $Q_4 - M$ is $P_{ux} + P_{x_Ry_R} + P_{yv}$ where P_{ux} and P_{yv} are disjoint subpaths of P_{uv}, assuming that x is closer to u than y on the path P_{uv}.

3.2.5: The remaining possibility; that is, M is as on Figure 24(2)-(4), the configuration (u, v, M_L) is bad, and $\{u, v\} \neq \{b, c\}$. Choose an edge $xy \in M_L$ such that the configurations $(u, v, M_L \setminus \{xy\})$ and (x_R, y_R, M_R) are good. Observe on Figure 24 and Figure 1 that such an edge exists.

Thus there is a Hamiltonian path $P_{x_Ry_R}$ of $Q_R - M_R$. Since the configuration (u, v, M_L) is bad but the configuration $(u, v, M_L \setminus \{xy\})$ is good, there is a Hamiltonian path P_{uv} that contains edge xy and faults $M_L \setminus \{xy\}$. The desired Hamiltonian path of $Q_4 - M$ is $P_{ux} + P_{x_Ry_R} + P_{yv}$ where P_{ux} and P_{yv} are disjoint subpaths of P_{uv}, assuming that x is closer to u than y on the path P_{uv}. This completes the proof of the lemma. □

8 Gray codes in hypercubes of higher dimensions

In this section we use $\mathrm{AL}(u, v, M, Q_n)$ to denote the fact that a matching M contains an almost-layer in Q_n with a free pair $\{u, v\}$. Recall that if $\mathrm{AL}(u, v, M, Q_n)$ holds, then there exists d such that M^d consists of all edges of dimension d except uu^d and vv^d. The following simple observation is used later.

Proposition 8.1. *Let M be a matching which contains an almost-layer A with a free pair $\{u, v\}$ in Q_n, $n \geq 4$.*

(*i*) *If* $\mathrm{AL}(u, w, M, Q_n)$ *holds, then* $v = w$.

(*ii*) *If at least one of vertices* x, y *is incident with an edge of* A*, then* $\mathrm{AL}(x, y, M, Q_n)$ *does not hold.*

Note that the assumption $n \geq 4$ is for part (*ii*) essential, as it fails to hold in case that M is the perfect matching of Q_3 depicted on Figure 2(10). Now we are ready to resolve the case of matchings not containing a half-layer.

Lemma 8.2. *Let* $u, v \in V(Q_n)$*,* $n \geq 4$*, such that* $p(u) \neq p(v)$ *and let* M *be a matching in* Q_n *containing no half-layer. Then* $Q_n - M$ *contains a Hamiltonian path between* u *and* v *unless* M *is an almost-layer in* Q_n *with a free pair* $\{u, v\}$*.*

Proof. We argue by induction on n. Since the case $n = 4$ is settled by Lemma 7.3, assume that $n > 4$. If M is aligned, the statement follows from Lemma 5.3. If M is unaligned, then by Lemma 4.2, there is $d \in [n]$ such that M_i^d contains no half-layer in Q_i^d for both $i \in \{L, R\}$. Without a loss of generality assume that $u \in V(Q_L^d)$. To obtain the desired Hamiltonian path P_{uv} of $Q_n - M$, consider the following two cases:

Case 1: $v \in Q_R^d$. Since M does not contain a half-layer, there is a vertex $w \in V(Q_L^d)$ such that

$$p(u) \neq p(w) \quad \text{and} \quad ww^d \notin M. \tag{1}$$

Notice that also $p(w^d) \neq p(v)$. We claim that w can be chosen so that neither $\mathrm{AL}(u, w, M_L^d, Q_L^d)$ nor $\mathrm{AL}(w^d, v, M_R^d, Q_R^d)$ holds. Indeed, if M_L^d contains an almost-layer A in Q_L^d with a free pair $\{u, w\}$, then each edge of A contains one vertex that can play the role of w in (1). Hence there are still $|A| = 2^{n-2} - 2 > 2$ choices for w, and M_R^d may contain an almost-layer in Q_R^d with a free pair $\{w^d, v\}$ for no more than one of them by Proposition 8.1(*i*).

It remains to apply the induction to obtain Hamiltonian paths P_{uw} and $P_{w^d v}$ of $Q_L^d - M_L^d$ and $Q_R^d - M_R^d$, respectively, and finally put $P_{uv} := P_{uw} + P_{w^d v}$ to conclude this case.

Case 2: $v \in V(Q_L^d)$. Consider the following subcases.

Subcase 2.1: $\mathrm{AL}(u, v, M_L^d, Q_L^d)$ *holds.* Since M_L^d contains an almost-layer A with $2^{n-2} - 2 > 3$ edges, Proposition 8.1(*ii*) implies that there must be an edge $xy \in A$ such that M_R^d contains no almost-layer in Q_R^d with a free pair $\{x^d, y^d\}$. Thus, by the induction hypothesis, there are Hamiltonian paths P_{uv} and $P_{x^d y^d}$ of $Q_L^d - (M_L^d \setminus \{xy\})$ and $Q_R^d - M_R^d$, respectively. Note that by Lemma 5.1, the path P_{uv} passes through the edge xy, and hence there are subpaths P_{ux} and P_{yv} of P_{uv} such that $P_{uv} = P_{ux} + P_{yv}$. It remains to put $P_{uv} := P_{ux} + P_{x^d y^d} + P_{yv}$ to complete this subcase.

Subcase 2.2: $\mathrm{AL}(u, v, M_L^d, Q_L^d)$ *does not hold.* First apply the induction hypothesis to obtain a Hamiltonian path P_{uv}^L of $Q_L^d - M_L^d$. Next, distinguish the following two possibilities.

Subcase 2.2.1: $\{xx^d, yy^d\} \cap M = \emptyset$ *for some pair* x, y *of neighbors on* P_{uv}^L. We claim that vertices x and y can be chosen so that $\mathrm{AL}(x^d, y^d, M_R^d, Q_R^d)$ does not hold. To verify the claim, suppose, by way of contradiction, that M_R^d contains an almost-layer A in Q_R^d with a free pair $\{x^d, y^d\}$. Denote by S the set of vertices z of Q_L^d such that z^d is incident with an edge of A. Then

$$|S| = 2|A| = 2^{n-1} - 4 > 2^{n-2} = |V(Q_L^d)|/2 = |V(P_{uv}^L)|/2$$

and hence there exist $x, y \in S$ that are neighbors on P_{uv}^L. Proposition 8.1(*ii*) then implies that $\mathrm{AL}(x^d, y^d, M_R^d, Q_R^d)$ does not hold.

As in Subcase 2.1, it remains to apply the induction to obtain a Hamiltonian path $P_{x^d y^d}$ of $Q_R^d - M_R^d$ and put $P_{uv} := P_{ux} + P_{x^d y^d} + P_{yv}$, where P_{ux} and P_{yv} are the subpaths of P_{uv}^L satisfying $P_{uv}^L = P_{ux} + P_{yv}$.

Subcase 2.2.2: $\{xx^d, yy^d\} \cap M \neq \emptyset$ *for each pair* x, y *of neighbors on* P_{uv}^L. By Lemma 4.2 there exists another splitting dimension $d' \neq d$ and we claim that Subcase 2.2.2 for d' cannot occur. To verify the claim, first observe that

$$|M^d| \geq |V(P_{uv}^L)|/2 = 2^{n-2}. \tag{2}$$

If this inequality is strict, then $|M| \leq 2^{n-1}$ implies that $|M^{d'}| < 2^{n-2}$, contrary to (2), and hence the claim holds. Thus we can assume that

$$|M^d| = |M^{d'}| = 2^{n-2}. \tag{3}$$

This means that $|M^d| = |V(P_{uv}^L)|/2$, which together with the assumption of this subcase implies that $|\{xx^d, yy^d\} \cap M^d| = 1$ for each pair x, y of neighbors on P_{uv}^L. Since the parity of consecutive vertices on P_{uv}^L alternates, it follows that if $uu^d \in M^d$, then $xx^d \in M^d$ for each vertex $x \in V(P_{uv}^L)$ with $p(u) = p(x)$. However, this would mean that M contains a half-layer, contrary to our assumption. Since the same reasons apply to the endvertex v, we conclude that neither u nor v are incident with an edge of M^d. Applying the same argument to the dimension d', it follows that

$$|M^d| + |M^{d'}| \leq |V(Q_n) \setminus \{u, v\}|/2 = 2^{n-1} - 1,$$

contrary to (3). □

9 The main results

Now we are ready to formulate the necessary and sufficient conditions for (cyclic) Gray codes faulting a given matching. Note that the conditions are expressed in terms of half-layers and almost-layers, which were defined in Sections 4 and 5, respectively.

Theorem 9.1. *Let M be a matching, and let u and v be vertices of Q_n, $n \geq 4$. Then Q_n has a Gray code between u and v faulting M if and only if M contains no odd half-layer for u or v and no almost layer with a free pair $\{u, v\}$.*

Proof. Recall that (u, v, M) is an admissible configuration if and only if the left side of the equality holds, by Proposition 5.2. The necessity is settled by Lemma 5.1. The sufficiency follows from Lemma 5.3 in case M contains a half-layer, and from Lemma 8.2 otherwise. □

It should be noted that for $n = 3$ the problem is fully described by Proposition 6.1 in Section 6, and for $n \leq 2$ it is trivial.

Corollary 9.2. *Let M be a matching in Q_n, $n \geq 4$. Then Q_n has a cyclic Gray code faulting M if and only if M does not contain a half-layer.*

Proof. Observe that the following two equivalences hold.

- $Q_n - M$ contains a Hamiltonian cycle $\iff$ $Q_n - M$ contains a Hamiltonian path P_{uv} for some edge $uv \in E(Q_n)$;
- there exists an admissible configuration (u, v, M) for some edge $uv \in E(Q_n)$ $\iff$ M does not contain a half-layer.

The corollary now follows from Theorem 9.1. □

Corollary 9.3. *Let M be a perfect matching in Q_n. Then the following three statements are equivalent:*

(*i*) *Q_n has a Gray code faulting M;*

(*ii*) *Q_n has a cyclic Gray code faulting M;*

(*iii*) *$Q_n - M$ is connected.*

Proof. First observe that clearly $(ii) \implies (i)$ and $(i) \implies (iii)$. Hence it only remains to prove that $(iii) \implies (ii)$. To that end, notice that for any perfect matching M of Q_n,

M contains a half-layer in Q_n $\iff$
M forms a layer of Q_n $\iff$
$Q_n - M$ is disconnected.

The validity of the implication $(iii) \implies (ii)$ for $n \geq 4$ now follows from Corollary 9.2, the case $n = 3$ may be verified by inspection of parts (9) and (10) of Figure 2, and the case $n \leq 2$ is trivial. □

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