

Ramsey theorem for designs

Jan Hubička^{1,2} Jaroslav Nešetřil^{1,3}

*Computer Science Institute of Charles University (IUUK)
Charles University
Prague, Czech Republic*

Abstract

We prove that for any choice of parameters k, t, λ the class of all finite ordered designs with parameters k, t, λ is a Ramsey class.

Keywords: Ramsey class, homogeneous structure, design, Steiner system

1 Introduction

We prove that for every choice of parameters $2 \leq t \leq k$ and $1 \leq \lambda$ the class $\overrightarrow{PD}_{kt\lambda}$ of linearly ordered partial designs with parameters k, t, λ is a Ramsey class. Thus, together with the recent spectacular results of Keevash [11], one obtains that the class of linearly ordered designs $\overrightarrow{D}_{kt\lambda}$ is a Ramsey class.

This paper involves three seemingly unrelated subjects: block designs, model theory and structural Ramsey theory. The generality is an important issue and in such context our main result can be formulated as follows:

¹ Supported by grant ERC-CZ LL-1201 of the Czech Ministry of Education and CE-ITI P202/12/G061 of GAČR.

² Email: hubicka@iuuk.mff.cuni.cz

³ Email: nesetril@iuuk.mff.cuni.cz

Theorem 1.1 *For any choice of parameters k, t, λ the class $\vec{\mathcal{D}}_{kt\lambda}$ of all finite ordered designs with parameters k, t, λ is a Ramsey class.*

For the proof we have to find the right degree of abstraction which will be introduced in the next three sub-sections together with all relevant notions. Strong structural Ramsey theorem (proved in [8,5]) plays the key role.

1.1 Designs

A (k, t, λ) -*design* ($k \geq t, \lambda$ all positive integers) is a finite hypergraph (X, R) where R is a set of k -subsets of X with property that any t -subset of X is contained in exactly λ elements of R . More formally we have $R \subseteq \binom{X}{k}$ and $|\{M \in R : T \subseteq M\}| = \lambda$ for any $T \in \binom{X}{t}$ (as usual in Ramsey context we denote by $\binom{X}{k}$ the set of all k -subsets of X consisting of k elements). A *partial* (k, t, λ) -*design* is hypergraph (X, R) where every t -subset is in at most λ elements of R .

Designs form a classical area of combinatorics as well as of mathematical statistics (design of experiments). Particularly Keevash [11,9], extending another spectacular result in the area [14,15,16], recently showed the following:

Theorem 1.2 (Keevash theorem [11]) *For every choice of parameters k, t, λ there exists (k, t, λ) -design on every sufficiently large set satisfying a well known divisibility condition. Also any partial (k, t, λ) -design can be completed to a (k, t, λ) -design.*

1.2 Models

Let $L = L_{\mathcal{R}} \cup L_{\mathcal{F}}$ be a language involving relational symbols $R \in L_{\mathcal{R}}$ and function symbols $F \in L_{\mathcal{F}}$ each having associated positive integers called *arity* and denoted by $a(R)$ for relations and *domain arity*, $d(F)$, *range arity*, $r(F)$, for functions. An L -*structure* \mathbf{A} is a structure with *vertex set* A , functions $F_{\mathbf{A}} : \text{Dom}(F_{\mathbf{A}}) \rightarrow \binom{A}{r(F)}$, $\text{Dom}(F_{\mathbf{A}}) \subseteq A^{d(F)}$ for $F \in L_{\mathcal{F}}$ and relations $R_{\mathbf{A}} \subseteq A^{a(R)}$ for $R \in L_{\mathcal{R}}$. $\text{Dom}(F_{\mathbf{A}})$ is called the *domain* of function F in \mathbf{A} . Notice that the domain is set of ordered $d(F)$ -tuples while the range is set of unordered $r(F)$ -tuples. Symmetry in ranges permits explicit description of algebraic closures in the Fraïssé limits without changing the automorphism group (c.f. [4]). It also simplifies some of the notation bellow.

The language L is usually fixed and understood from the context. If set A is finite we call \mathbf{A} *finite structure* (in most of this paper all structures

are finite). If language L contains no function symbols, we call L *relational language* and every L -structure is also called *relational L -structure*.

The notion of embeddings, isomorphism, homomorphisms and free amalgamation are natural generalisation of the corresponding notions on relational structures and are formally introduced in Section 2. Considering function symbols has important consequences to what we consider as a substructure: An L -structure \mathbf{A} is a *substructure* of \mathbf{B} if $A \subseteq B$ and all relations and functions of \mathbf{B} restricted to A are precisely relations and functions of A . In particular a t -tuple \mathbf{t} of vertices of A is in $\text{Dom}(F_{\mathbf{A}})$ if and only if it is also in $\text{Dom}(F_{\mathbf{B}})$ and $F_{\mathbf{A}}(\mathbf{t}) = F_{\mathbf{B}}(\mathbf{t})$. This implies the fact that \mathbf{B} does not induce a substructure on every subset of B (but only on “closed” sets, to be defined later).

In our setting (k, t, λ) -design corresponds to a particular L -structure with $L_{\mathcal{F}} = \emptyset$, $L_{\mathcal{R}} = \{R\}$, $a(R) = k$. However designs satisfy some further properties and it is essential for our argument that we introduce and use function symbols.

1.3 Ramsey classes

For structures \mathbf{A}, \mathbf{B} denote by $\binom{\mathbf{B}}{\mathbf{A}}$ the set of all sub-structures of \mathbf{B} , which are isomorphic to \mathbf{A} . Using this notation the definition of a Ramsey class gets the following form: A class \mathcal{C} is a *Ramsey class* if for every two objects \mathbf{A} and \mathbf{B} in \mathcal{C} and for every positive integer k there exists a structure \mathbf{C} in \mathcal{C} such that the following holds: For every partition $\binom{\mathbf{C}}{\mathbf{A}}$ into k classes there exists an $\tilde{\mathbf{B}} \in \binom{\mathbf{C}}{\mathbf{B}}$ such that $\binom{\tilde{\mathbf{B}}}{\mathbf{A}}$ belongs to one class of the partition. It is usual to shorten the last part of the definition to $\mathbf{C} \rightarrow (\mathbf{B})_k^{\mathbf{A}}$.

We are motivated by the following, now classical, result.

Theorem 1.3 (Nešetřil-Rödl theorem [13]) *Let \mathbf{A} and \mathbf{B} be a linearly ordered hypergraphs, then there exists a linearly ordered hypergraph \mathbf{C} such that $\mathbf{C} \rightarrow (\mathbf{B})_2^{\mathbf{A}}$.*

Moreover, if \mathbf{A} and \mathbf{B} do not contain an irreducible hypergraph \mathbf{F} then \mathbf{C} may be chosen with the same property (a hypergraph F is irreducible if every pair of its elements is contained in an edge of F).

Given language L , denote by \overrightarrow{L} language L extended by one binary relation \leq . Given L -structure \mathbf{A} the *ordering of \mathbf{A}* is \overrightarrow{L} -structure extending \mathbf{A} by arbitrary linear ordering of vertices represented by $\leq_{\mathbf{A}}$. We denote such ordered \mathbf{A} as $\overrightarrow{\mathbf{A}}$. Given class \mathcal{K} of L -structures denote by $\overrightarrow{\mathcal{K}}$ the class of all orderings of structures in \mathcal{K} i.e. $\overrightarrow{\mathcal{K}}$ is a class of \overrightarrow{L} -structures where \leq is a linear order. We sometimes say that $\overrightarrow{\mathcal{K}}$ arises by the *free orderings* of structures in \mathcal{K} . For purposes of this paper Theorem 1.3 can now be re-formulated using notions

of Fraïssé theory (which will be briefly introduced in Section 2) as follows:

Theorem 1.4 (Ramsey theorem for free amalgamation classes) *Let L be a relational language, \mathcal{K} be a free amalgamation class of relational L -structures. Then $\overrightarrow{\mathcal{K}}$ is a Ramsey class.*

Connection of Ramsey classes and extremely amenable groups [10] motivated a systematic search for new examples of Ramsey classes. It became apparent that it is important to consider structures with both relations and functions or, equivalently, classes of structures with “strong embeddings”. This led to [8] which provides a sufficient structural condition for a subclass of a Ramsey class to be Ramsey and generalises this approach also to classes with function symbols representing closures. Comparing the two main results of [8] (Theorem 2.1 for classes without closures and Theorem 2.2 for classes with closures) it is clear that considering closures leads to many technical difficulties. In fact, a recent example given in [4] shows that there is no direct analogy of Theorem 1.4 to free amalgamation classes with closures. Perhaps surprisingly, one can prove that if closures are explicitly represented by means of partial functions, such statement is true. More precisely the following we proved in [5, Theorem 1.3] as a more streamlined version of Theorem 2.2 of [8]:

Theorem 1.5 *Let L be a language (involving relational symbols and partial functions), \mathcal{K} be a free amalgamation class of L -structures. Then $\overrightarrow{\mathcal{K}}$ is a Ramsey class.*

It appears (see [5]) that many natural classes may be interpreted as free amalgamation classes and consequently Theorem 1.5 yields uniform proofs of Ramsey property of some recently discovered Ramsey classes (such as ordered partial Steiner systems (i.e. $(k, t, 1)$ -designs) [1], bowtie-free graphs [7], bouquet-free graphs [3] and a Ramsey expansion of class 2-orientations of Hrushovski predimension construction [4]). In this paper we add to this list of applications of Theorem 1.5 yet another example from a very different area.

2 Preliminaries

We now review some standard model-theoretic notions (see e.g. [6]).

An *embedding* $f : \mathbf{A} \rightarrow \mathbf{B}$ is an injective mapping $f : A \rightarrow B$ satisfying for every $R \in L_{\mathcal{R}}$ and $F \in L_{\mathcal{F}}$:

- (i) $(x_1, x_2, \dots, x_{a(R)}) \in R_{\mathbf{A}} \iff (f(x_1), f(x_2), \dots, f(x_{a(R)})) \in R_{\mathbf{B}}$, and,
- (ii) $(x_1, \dots, x_{d(F)}) \in \text{Dom}(F_{\mathbf{A}}) \iff (f(x_1), \dots, f(x_{d(F)})) \in \text{Dom}(F_{\mathbf{B}})$ and $f(F_{\mathbf{A}}(x_1, c_2, \dots, x_{d(F)})) = F_{\mathbf{B}}(f(x_1), f(x_2), \dots, f(x_{d(F)}))$.

If f is an embedding which is an inclusion then \mathbf{A} is a *substructure* of \mathbf{B} . For an embedding $f : \mathbf{A} \rightarrow \mathbf{B}$ we say that \mathbf{A} is *isomorphic* to $f(\mathbf{A})$ and $f(\mathbf{A})$ is also called a *copy* of \mathbf{A} in \mathbf{B} . Thus $\binom{\mathbf{B}}{\mathbf{A}}$ is defined as the set of all copies of \mathbf{A} in \mathbf{B} . Given $\mathbf{A} \in \mathcal{K}$ and $B \subset A$, the *closure of B in \mathbf{A}* is the smallest substructure of \mathbf{A} containing B .

Let \mathbf{A} , \mathbf{B}_1 and \mathbf{B}_2 be structures with α_1 an embedding of \mathbf{A} into \mathbf{B}_1 and α_2 an embedding of \mathbf{A} into \mathbf{B}_2 , then every structure \mathbf{C} together with embeddings $\beta_1 : \mathbf{B}_1 \rightarrow \mathbf{C}$ and $\beta_2 : \mathbf{B}_2 \rightarrow \mathbf{C}$ satisfying $\beta_1 \circ \alpha_1 = \beta_2 \circ \alpha_2$ is called an *amalgamation of \mathbf{B}_1 and \mathbf{B}_2 over \mathbf{A} with respect to α_1 and α_2* . We will call \mathbf{C} simply an *amalgamation* of \mathbf{B}_1 and \mathbf{B}_2 over \mathbf{A} (as in the most cases α_1, α_2 and β_1, β_2 can be chosen to be inclusion embeddings).

Amalgamation is *free* if $\beta_1(x_1) = \beta_2(x_2)$ if and only if $x_1 \in \alpha_1(A)$ and $x_2 \in \alpha_2(A)$ and there are no tuples in any relations of \mathbf{C} and $\text{Dom}(F_{\mathbf{C}})$, $F \in L_{\mathcal{F}}$, using both vertices of $\beta_1(B_1 \setminus \alpha_1(A))$ and $\beta_2(B_2 \setminus \alpha_2(A))$. An *amalgamation class* is a class \mathcal{K} of finite structures satisfying the following three conditions:

- (i) *Hereditary property*: For every $\mathbf{A} \in \mathcal{K}$ and a substructure \mathbf{B} of \mathbf{A} we have $\mathbf{B} \in \mathcal{K}$;
- (ii) *Joint embedding property*: For every $\mathbf{A}, \mathbf{B} \in \mathcal{K}$ there exists $\mathbf{C} \in \mathcal{K}$ such that \mathbf{C} contains both \mathbf{A} and \mathbf{B} as substructures;
- (iii) *Amalgamation property*: For $\mathbf{A}, \mathbf{B}_1, \mathbf{B}_2 \in \mathcal{K}$ and α_1 embedding of \mathbf{A} into \mathbf{B}_1 , α_2 embedding of \mathbf{A} into \mathbf{B}_2 , there is $\mathbf{C} \in \mathcal{K}$ which is an amalgamation of \mathbf{B}_1 and \mathbf{B}_2 over \mathbf{A} with respect to α_1 and α_2 .

If the \mathbf{C} in the amalgamation property can always be chosen as the free amalgamation, then \mathcal{K} is *free amalgamation class*.

This explains all the notions in our Theorem 1.5. In the next section we apply this result to designs.

3 Main result

To deal with partial (k, t, λ) -designs in order to apply Theorem 1.5 we need a particular encoding. Our language is denoted by $L = (L_{\mathcal{R}}, L_{\mathcal{F}})$. The relational part $L_{\mathcal{R}}$ consists from relational symbol R of arity k . We put $K = (k - t)\lambda + t$. The functional language $L_{\mathcal{F}}$ consists from symbol F^k, F^{k+1}, \dots, F^K all with domain arity $d(F^\ell) = t$ and range arity $r(F^\ell) = \ell$, $\ell = k, k + 1, \dots, K$.

Denote by $\text{Str}(L)$ the class of all L -structures (i.e. models of the language L). Within this class $\text{Str}(L)$ we define a subclass $\mathcal{PD}_{kt\lambda}$ of all structures $\mathbf{A} = (A, R_{\mathbf{A}}, (F_{\mathbf{A}}^\ell : \ell = k, k + 1, \dots, K))$ which satisfy

- (i) for every $(v_1, v_2, \dots, v_k) \in R_{\mathbf{A}}$ it holds that $v_i \neq v_j$, $1 \leq i < j \leq k$;
- (ii) if $(v_1, v_2, \dots, v_k) \in R_{\mathbf{A}}$ then $(v_{\pi(1)}, v_{\pi(2)}, \dots, v_{\pi(k)}) \in R_{\mathbf{A}}$ for any permutation π ;
- (iii) $(A, \tilde{R}_{\mathbf{A}})$ is partial (k, t, λ) -design where $\tilde{R}_{\mathbf{A}} = \{\{v_1, v_2, \dots, v_k\} : (v_1, v_2, \dots, v_k) \in R_{\mathbf{A}}\}$;
- (iv) $(v_1, v_2, \dots, v_t) \in \text{Dom}(F_{\mathbf{A}}^\ell)$ if and only if $|N_{\mathbf{A}}(v_1, v_2, \dots, v_t)| = \ell$ and $F_{\mathbf{A}}^\ell(v_1, v_2, \dots, v_t) = N_{\mathbf{A}}(v_1, v_2, \dots, v_t)$ where $N_{\mathbf{A}}(v_1, v_2, \dots, v_t)$ is the *neighbourhood* of set $\{v_1, v_2, \dots, v_t\}$ — that is the set of all vertices v such that there exists $M \in \tilde{R}_{\mathbf{A}}$ containing each of vertices v, v_1, v_2, \dots, v_t .

The embeddings in $\mathcal{PD}_{kt\lambda}$ are inherited from $\text{Str}(L)$. Let us explicitly formulate their form:

For $\mathbf{A} = (A, R_{\mathbf{A}}, (F_{\mathbf{A}}^\ell : \ell = k, k+1, \dots, K))$, $\mathbf{B} = (B, R_{\mathbf{B}}, (F_{\mathbf{B}}^\ell : \ell = k, k+1, \dots, K))$ injective mapping $f : A \rightarrow B$ is an embedding of \mathbf{A} into \mathbf{B} if it satisfies the following conditions:

- (i) $(x_1, x_2, \dots, x_k) \in R_{\mathbf{A}}$ if and only if $(f(x_1), f(x_2), \dots, f(x_k)) \in R_{\mathbf{B}}$ (i.e. f is embedding $(A, R_{\mathbf{A}})$ into $(B, R_{\mathbf{B}})$),
- (ii) for every ℓ , $k \leq \ell \leq K$, it satisfies $\{f(x) : x \in F_{\mathbf{A}}^\ell(x_1, x_2, \dots, x_t)\} = F_{\mathbf{B}}^\ell(\{f(x_1), f(x_2), \dots, f(x_t)\})$ whenever one side of this equation make sense.

Every ordered partial (k, t, λ) -design (X, R, \leq) may be interpreted in $\overrightarrow{\mathcal{PD}}_{kt\lambda}$ (the class of free orderings of $\mathcal{PD}_{kt\lambda}$) as the following L -structure: $\mathbf{A} = (A, R_{\mathbf{A}}, \leq_{\mathbf{A}}, (F_{\mathbf{A}}^\ell : k \leq \ell \leq K))$ where:

- (i) $A = X$,
- (ii) $R_{\mathbf{A}} = \{(v_1, v_2, \dots, v_k) : \{v_1, v_2, \dots, v_k\} \in R \text{ and } |\{v_1, v_2, \dots, v_k\}| = k\}$,
- (iii) $\leq_{\mathbf{A}} = \leq$.
- (iv) $F_{\mathbf{A}}^\ell(\mathbf{t})$ is defined for every t -tuple $\mathbf{t} = (t_1, t_2, \dots, t_t)$ without repeated vertices whenever $|\bigcup\{M : T \subseteq M \in R\}| = \ell$ and in this case $F_{\mathbf{A}}^\ell(\mathbf{t}) = \bigcup\{M : T \subseteq M \in R\} = N_{\mathbf{A}}(\mathbf{t})$ where $T = \{t_1, t_2, \dots, t_t\}$.

Clearly $\mathbf{A} \in \overrightarrow{\mathcal{PD}}_{kt\lambda}$ as it satisfies the above 4 conditions defining the class $\mathcal{PD}_{kt\lambda}$ and $\leq_{\mathbf{A}}$ is a linear order. Note also that this correspondence is 1-to-1 as every $\mathbf{A} \in \overrightarrow{\mathcal{PD}}_{kt\lambda}$ leads to an ordered partial (k, t, λ) -design $(A, \tilde{R}_{\mathbf{A}}, \leq_{\mathbf{A}})$.

The embeddings in $\mathcal{PD}_{kt\lambda}$ have the following meaning in the class of designs: $f : \mathbf{A} \rightarrow \mathbf{B}$ is an embedding if it satisfies (i), (ii) and the following:

- (iii') Every $M \in \tilde{R}_{\mathbf{B}} \setminus f(\tilde{R}_{\mathbf{A}})$ intersects the set $f(\mathbf{A})$ in at most $t-1$ elements (of course we have $f(A) = \{f(a) : a \in A\}$ and $f(\tilde{R}_{\mathbf{A}}) = \{f(M) : M \in \tilde{R}_{\mathbf{A}}\}$).

The equivalence (iii) and (iii') follows from our assumptions in definition of $PD_{kt\lambda}$. In this case we the set $f(A)$ is closed in \mathbf{B} . (Note that the condition (iii') is vacuous if \mathbf{A}, \mathbf{B} are (k, t, λ) -designs.)

The following is the main result of this note:

Theorem 3.1 *For any $k \geq t \geq 2$ and $\lambda \geq 1$ the class $\overrightarrow{PD}_{kt\lambda}$ is Ramsey.*

Proof. (sketch) We apply Theorem 1.5 to the class $\overrightarrow{PD}_{kt\lambda}$. Thus the only thing we have to check is the free amalgamation of $PD_{kt\lambda}$. Thus let $\mathbf{A}, \mathbf{B}_1, \mathbf{B}_2$ be structures in $PD_{kt\lambda}$, $\alpha_i : \mathbf{A} \rightarrow \mathbf{B}_i$ inclusion embeddings. Let $\mathbf{C} = (C, R_{\mathbf{C}}, (F_{\mathbf{C}}^{\ell} : k \leq \ell \leq K))$ be defined as follows: $(C, R_{\mathbf{C}})$ is the free amalgam of relational structures $(B_1, R_{\mathbf{B}_1})$ and $(B_2, R_{\mathbf{B}_2})$ over $(A, R_{\mathbf{A}})$. For $\ell = k, k+1, \dots, K$ we put $F_{\mathbf{C}}^{\ell}(T) = F_{\mathbf{B}_i}^{\ell}(T)$ whenever $F_{\mathbf{B}_i}^{\ell}(T)$ is defined (here we, without loss of generality, assume that embeddings β_1 and β_2 from the definition of amalgam are inclusions). Note that this definition is consistent as $\alpha_i(A)$ is a closed set in \mathbf{B}_i , $i = 1, 2$. Thus \mathbf{C} is a free amalgam of $\mathbf{A}, \mathbf{B}_1, \mathbf{B}_2$ and Theorem 1.5 applies. \square

This theorem together with Theorem 1.2 implies Theorem 1.1.

4 Remarks

1. The class $PD_{kt\lambda}$ is the age of an ultrahomogeneous Fraïssé limit $\mathbf{U}_{kt\lambda}$ which, by the Kechris, Pestov, Todorčević correspondence [10], is countable “geometry-like” structure with extremely amenable group of automorphisms and uniquely defined universal minimal flow (compare [2]).
2. Note that the essential feature of the above proof is generality of Theorem 1.5 and use of function (symbols) leading to the right definition of closed sets. It is easy to see that not closed subset do not form a Ramsey class (as, for example, one can distinguish subsets by their closures). Note also that by [5] the (k, t, λ) -designs have the ordering property, see also [12].
3. This proof and the relationship of designs and models has some further consequences and leads to interesting problems (such as the extension property for partial automorphisms (EPPA)), compare [4].

References

- [1] V. Bhat, J. Nešetřil, C. Reiher, and V. Rödl. A Ramsey class for Steiner systems. arXiv:1607.02792, 2016.

- [2] G. Cherlin, S. Shelah, and N. Shi. Universal graphs with forbidden subgraphs and algebraic closure. *Advances in Applied Mathematics*, 22(4):454–491, 1999.
- [3] G. Cherlin and L. Tallgren. Universal graphs with a forbidden near-path or 2-bouquet. *Journal of Graph Theory*, 56(1):41–63, 2007.
- [4] D. M. Evans, J. Hubička, and J. Nešetřil. Automorphism groups and Ramsey properties of sparse graphs. In preparation.
- [5] D. M. Evans, J. Hubička, and J. Nešetřil. Ramsey properties and extending partial automorphisms for classes of finite structures. In preparation.
- [6] W. Hodges. *Model theory*, volume 42. Cambridge University Press, 1993.
- [7] J. Hubička and J. Nešetřil. Bowtie-free graphs have a Ramsey lift. arXiv:1402.2700, 22 pages, 2014.
- [8] J. Hubička and J. Nešetřil. All those Ramsey classes (Ramsey classes with closures and forbidden homomorphisms). Submitted, arXiv:1606.07979, 58 pages, 2016.
- [9] G. Kalai. Designs exist! [after Peter Keevash]. *Séminaire BOURBAKI, 67ème année*, n° 1100:1–17, 2014–2015.
- [10] A. S. Kechris, V. G. Pestov, and S. Todorčević. Fraïssé limits, Ramsey theory, and topological dynamics of automorphism groups. *Geometric and Functional Analysis*, 15(1):106–189, 2005.
- [11] P. Keevash. The existence of designs. arXiv:1401.3665, 2014.
- [12] J. Nešetřil. Ramsey theory. In R. L. Graham, M. Grötschel, and L. Lovász, editors, *Handbook of Combinatorics*, pages 1331–1403. Elsevier, 1995.
- [13] J. Nešetřil and V. Rödl. The Ramsey property for graphs with forbidden complete subgraphs. *Journal of Combinatorial Theory, Series B*, 20(3):243–249, 1976.
- [14] R. M. Wilson. An existence theory for pairwise balanced designs I: Composition theorems and morphisms. *Journal of Combinatorial Theory, Series A*, 13(2):220–245, 1972.
- [15] R. M. Wilson. An existence theory for pairwise balanced designs II: The structure of PBD-closed sets and the existence conjectures. *Journal of Combinatorial Theory, Series A*, 13(2):246–273, 1972.
- [16] R. M. Wilson. An existence theory for pairwise balanced designs, III: Proof of the existence conjectures. *Journal of Combinatorial Theory, Series A*, 18(1):71–79, 1975.